

Bitcoin Privacy

João Pedro Monteiro

Overview

1. What is Bitcoin?
2. Privacy Issues
3. ZeroCoin
4. eZC: ZeroCoin Reloaded

What is Bitcoin?

“We're in the 21st Century and I can call someone in Indonesia,
see him on screen and talk to him for free...

“We're in the 21st Century and I can call someone in Indonesia,
see him on screen and talk to him for free...
... and yet I can't send him **1 cent.**”

(Wences Casares, Xapo CEO)



WikiLeaks ✓
@wikileaks

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WikiLeaks now accepts anonymous
Bitcoin donations on
1HB5XMLmzFVj8ALj6mfBsbifRoD4miY36v



RETWEETS
289

FAVORITES
38



4:12 PM - 14 Jun 2011

WikiLeaks

Addresses are identifiers which you use to send bitcoins to another person.

Summary

Address [1HB5XMLmzFVj8ALj6mfBsbifRoD4miY36v](#)

Hash 160 [b169f2b0b866db05900b93a5d76345f18d3afb24](#)

Tools [Taint Analysis](#) - [Related Tags](#) - [Unspent Outputs](#)

Transactions

No. Transactions 2931 

Total Received **\$ 955,004.61** 

Final Balance **\$ 143.31** 

[Request Payment](#)

[Donation Button](#)

PayPal stops payment delivery to Mega, citing 'business reasons'

By [Russell Brandom](#) on February 27, 2015 04:25 pm [✉ Email](#) [🐦 @russellbrandom](#)



Kim Dotcom 
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Let's give Bitcoin a boost :-) [#Mega](#)

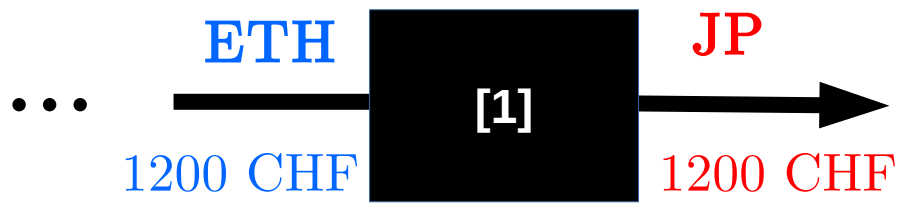
6:39 AM - 27 Feb 2015

7,636 RETWEETS **687** FAVORITES

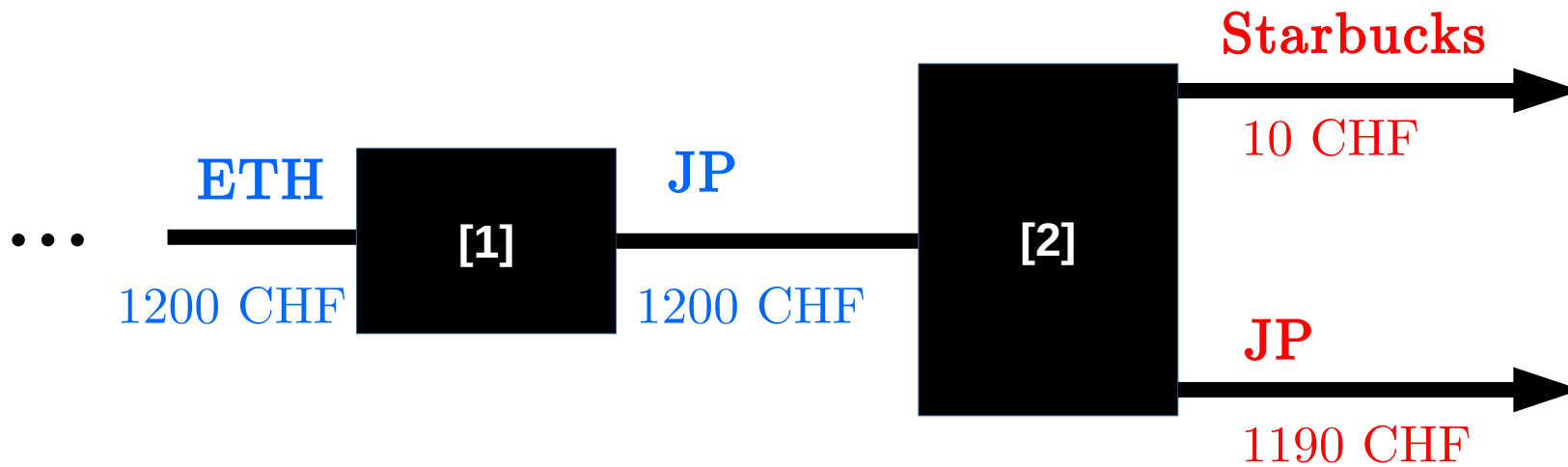




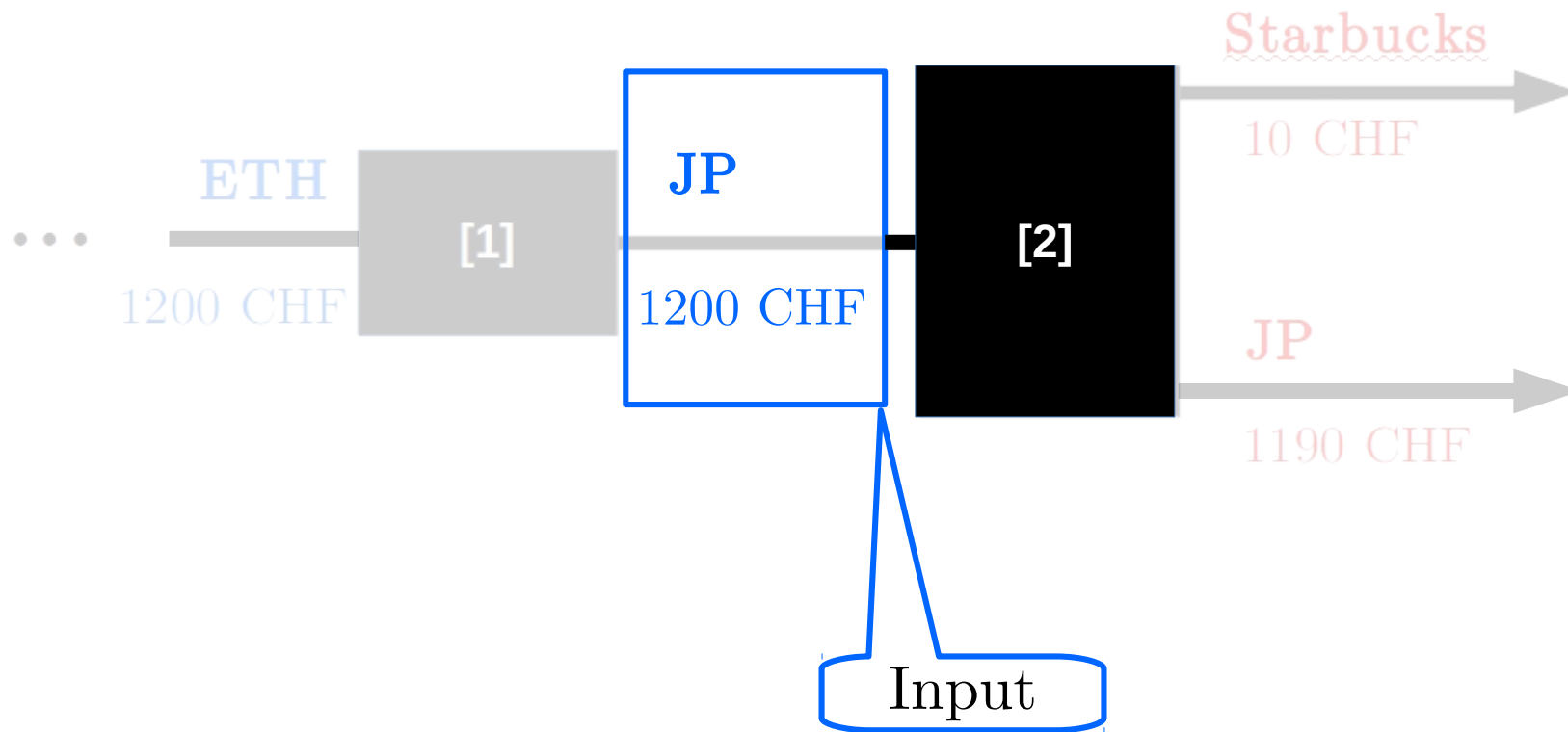
Transactions



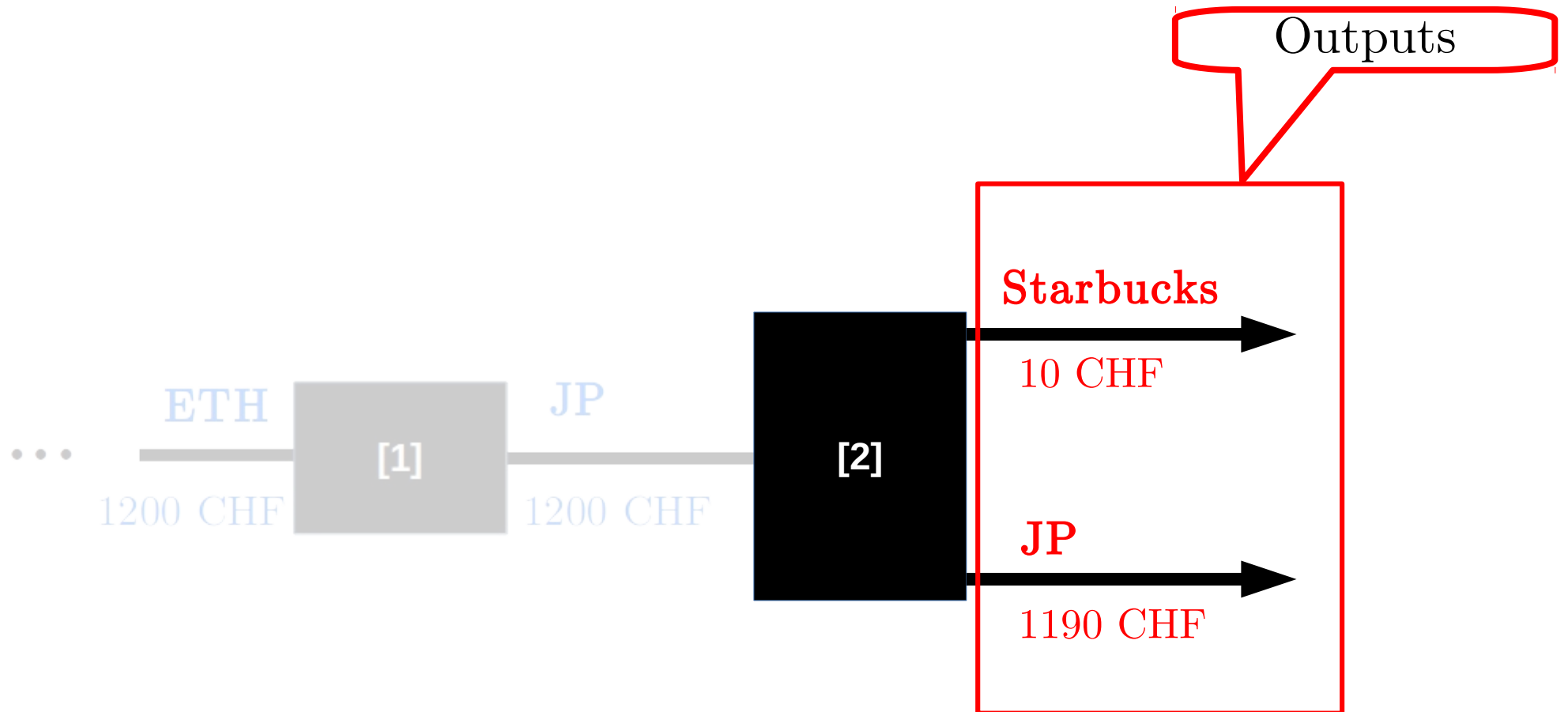
Transactions



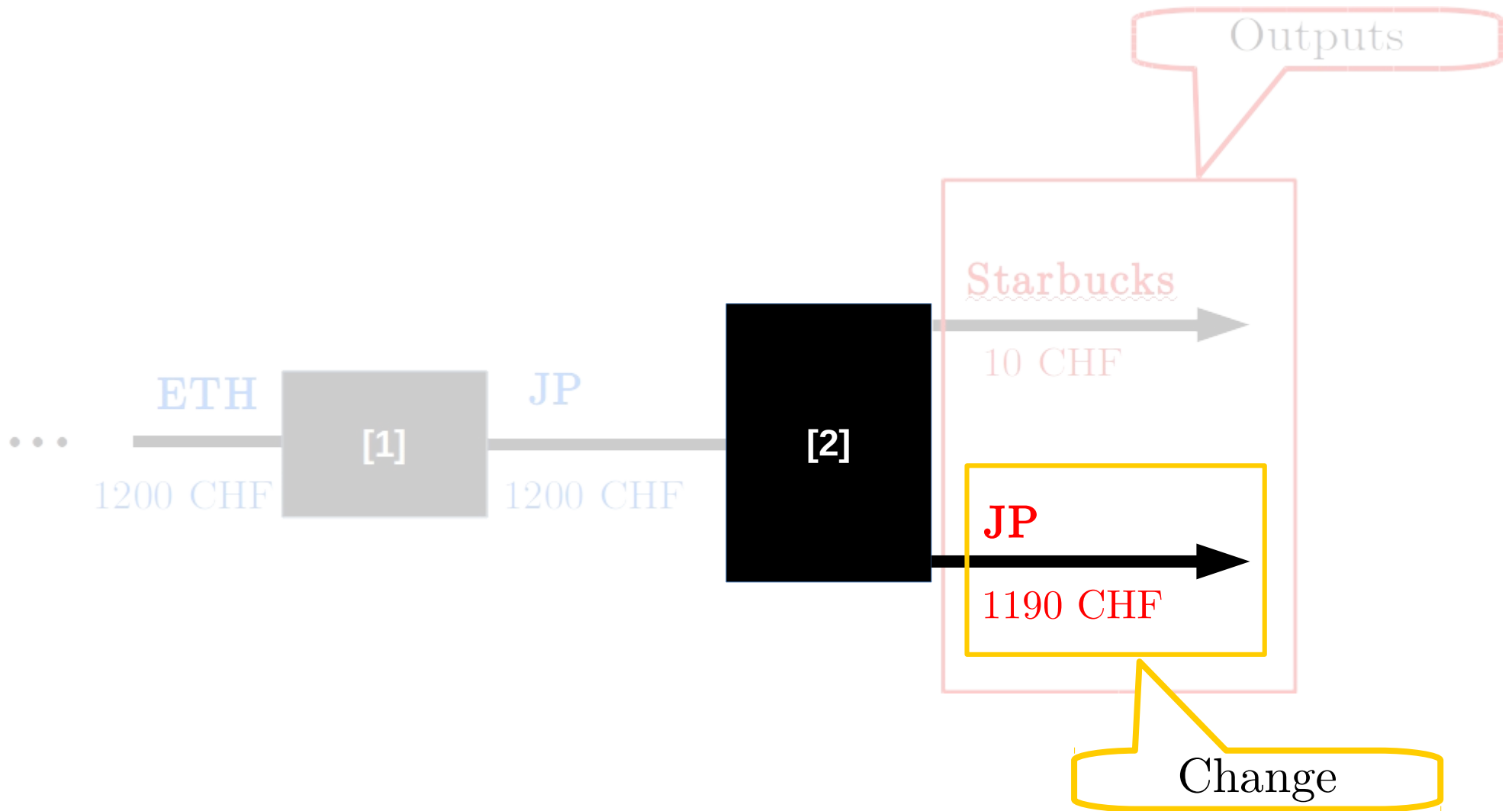
Transactions



Transactions



Transactions



Transactions

SAINT LUKE'S CREDIT UNION
4200 Pennsylvania, Suite 161
Kansas City, MO 64111

004485

18-81553010
001200155

REMITTER
ETH

DATE _____

PAY TO THE ORDER OF **J.P.**

1200 CHF

DOLLARS

VOID AFTER 90 DAYS

Eth zürich

THIS DOCUMENT HAS A MICRO-PRINT SIGNATURE LINE, WATERMARK AND A THERMOCHROMIC ICON; ABSENCE OF THESE FEATURES WILL INDICATE A COPY

VOID SAMPLE ONLY
NOT NEGOTIABLE

004485 3010813621 001200155

In a centralized system...



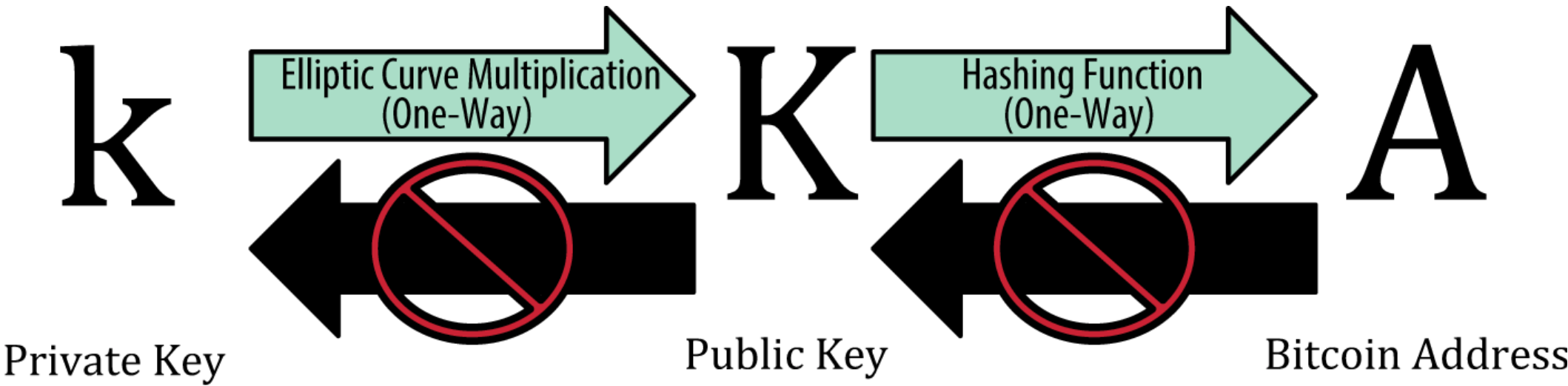
In a **decentralized** system...



In a **decentralized** system...



Key Triplets



Who can spend the created outputs?

Who can spend the created outputs?

Standard Transaction to Bitcoin address (pay-to-pubkey-hash)

```
scriptPubKey: OP_DUP OP_HASH160 <pubKeyHash> OP_EQUALVERIFY OP_CHECKSIG  
scriptSig: <sig> <pubKey>
```

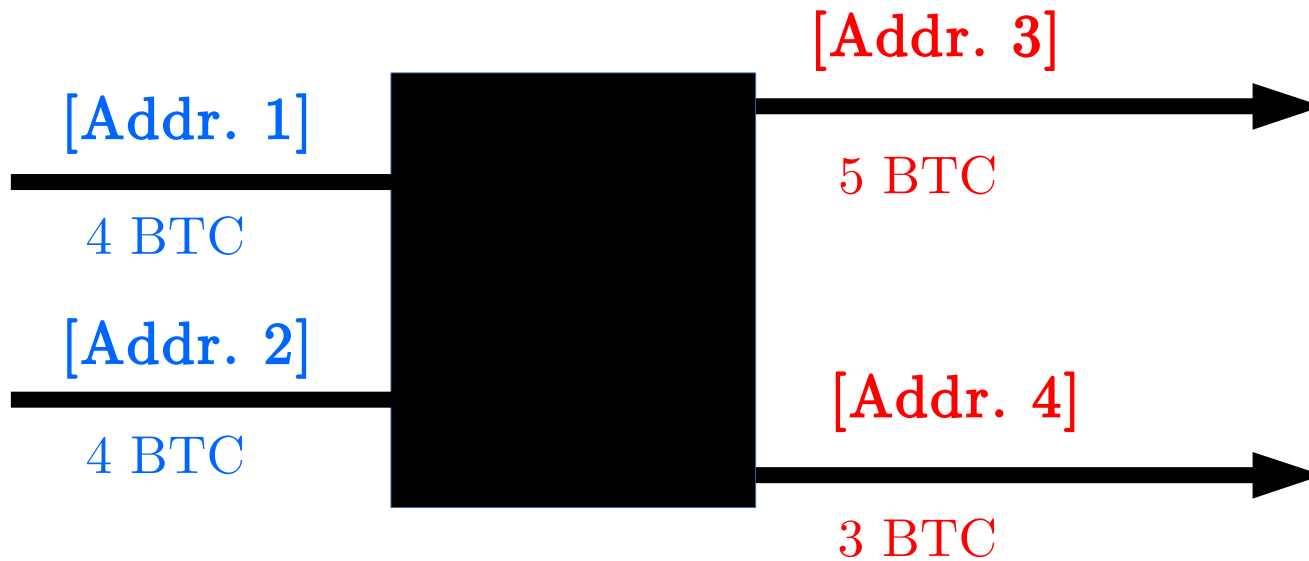
```
1
2
3 {
4   "txid" : "9ca8f969bd3ef5ec2a8685660fdbf7a8bd365524c2e1fc66c309acbae2c14ae3",
5   "version" : 1,
6   "locktime" : 0,
7   "vin" : [
8     {
9       "txid" : "d3c7e022ea80c4808e64dd0a1dba009f3eaae2318a4ece562f8ef815952717d7",
10      "vout" : 0,
11      "scriptSig" : {
12        "asm" :
13 "3045022100a4ebbeec83225dedead659bbde7da3d026c8b8e12e61a2df0dd0758e227383b302203301768ef878007e9ef7c304f70ffaf1f2c975b192d34c5b9b2ac1bd193dfba20104793ac8a58ea751f9710e39aad2e2
14 "483045022100a4ebbeec83225dedead659bbde7da3d026c8b8e12e61a2df0dd0758e227383b302203301768ef878007e9ef7c304f70ffaf1f2c975b192d34c5b9b2ac1bd193dfba2014104793ac8a58ea751f9710e39aa
15      },
16      "sequence" : 4294967295
17    }
18  ],
19  "vout" : [
20    {
21      "value" : 0.05000000,
22      "n" : 0,
23      "scriptPubKey" : {
24        "asm" : "OP_DUP OP_HASH160 07bdb518fa2e6089fd810235cf1100c9c13d1fd2 OP_EQUALVERIFY OP_CHECKSIG",
25        "hex" : "76a91407bdb518fa2e6089fd810235cf1100c9c13d1fd288ac",
26        "reqSigs" : 1,
27        "type" : "pubkeyhash",
28        "addresses" : [
29          "1hVzSofGwT8cjb8JU7nBsCSfEVQX5u9CL"
30        ]
31      }
32    },
33    {
34      "value" : 1.03362847,
35      "n" : 1,
36      "scriptPubKey" : {
37        "asm" : "OP_DUP OP_HASH160 107b7086b31518935c8d28703d66d09b36231343 OP_EQUALVERIFY OP_CHECKSIG",
38        "hex" : "76a914107b7086b31518935c8d28703d66d09b3623134388ac",
39        "reqSigs" : 1,
40        "type" : "pubkeyhash",
41        "addresses" : [
42          "12W9goQ3P7Waw5JH8fRVs1e2rVAKoGnvoy"
43        ]
44      }
45    }
46  ]
}
```

Privacy Issues with Bitcoin

**NOT SURE IF BITCOIN WAS CREATED
AGAINST THE NSA**

OR BY THE NSA

Isn't Bitcoin anonymous?!



Isn't Bitcoin anonymous?!

7e91e648710a99017d500c3953bb4e5dd684c97b235e50c8c52b3ec037e6bdcf

(Fee: 0.00 BTC - Size: 306 bytes) 2011-06-28 21:06:01

13vFf3MZKxSA3Q9e14c8xUXbMPhQn1wCgq (50 BTC - Output)

17zeTMh8xXeXXjZnbULXV3g3t3f7pftnEh (50 BTC - Output)



WikiLeaks [🔗](#) - (Spent)

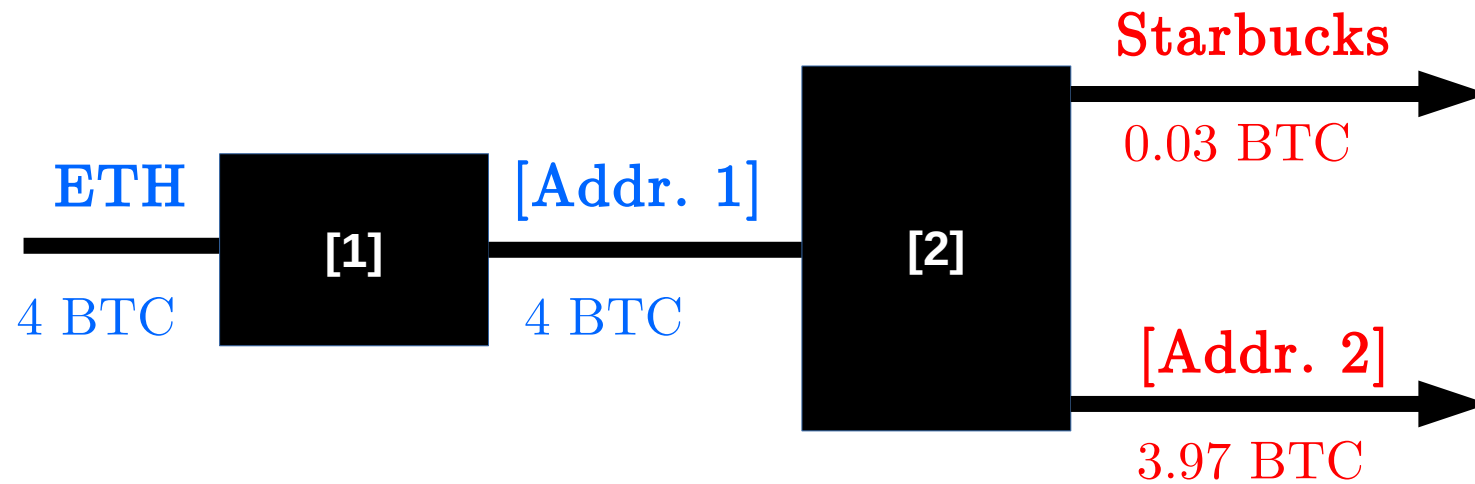
90 BTC

1Ne2mvY6Du6kQjjwi6tnHy2nLZup2n8Kgv - (Spent)

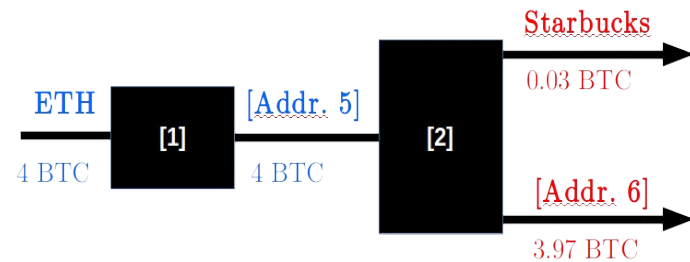
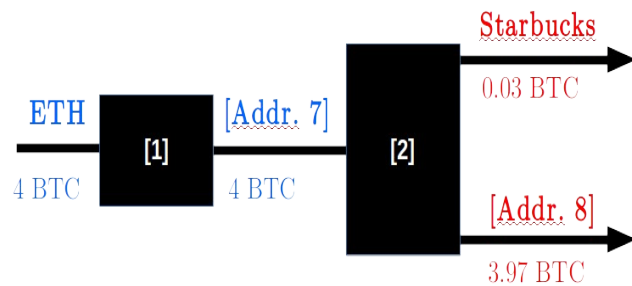
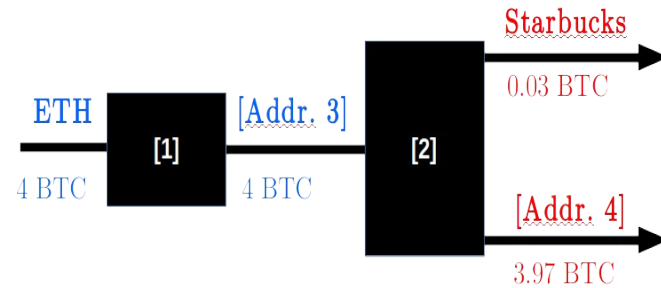
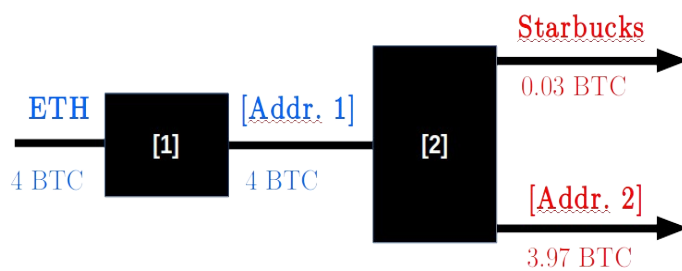
10 BTC

90 BTC

Isn't Bitcoin anonymous?!



Isn't Bitcoin anonymous?!



Why Anonymity?

"What we used to call liberty and freedom
we now call privacy."

(Jacob Appelbaum)



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1HB5XMLmzFVj8ALj6mfBsbifRoD4miY36v



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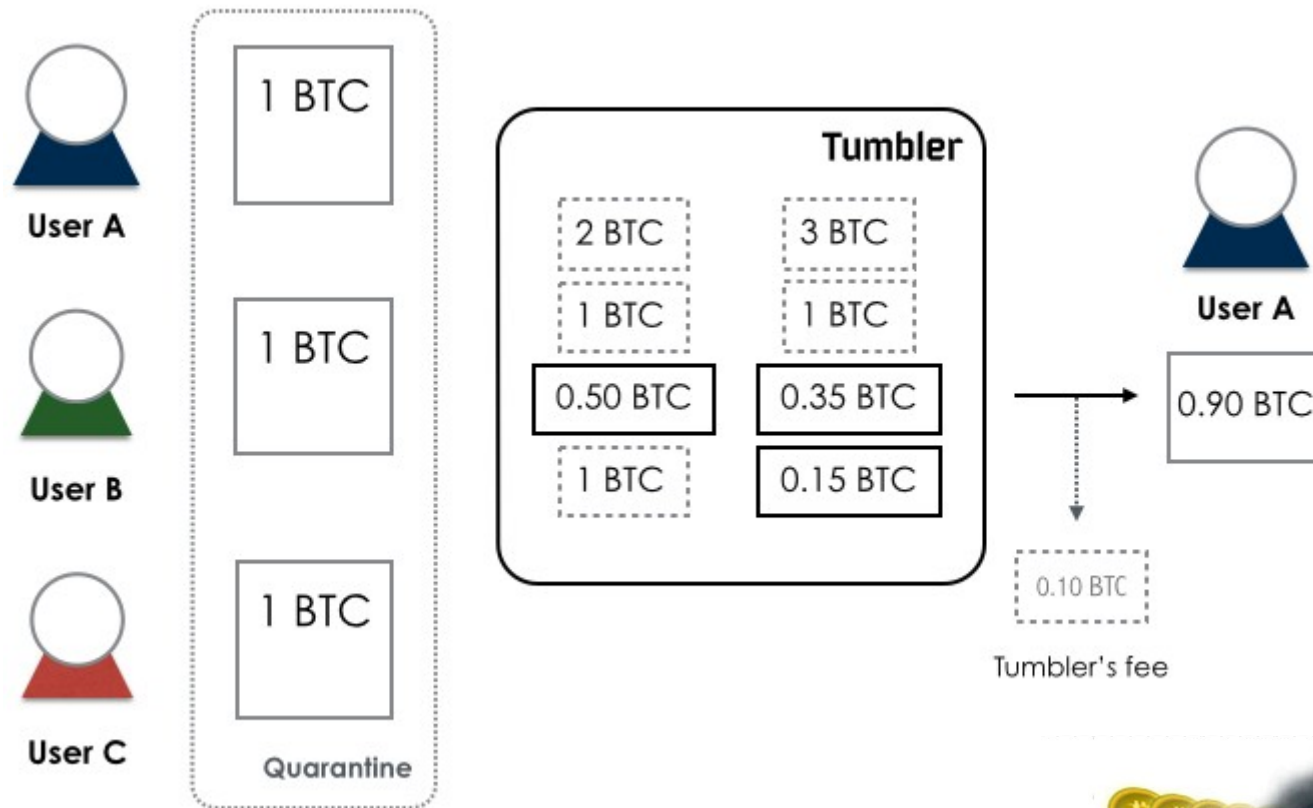
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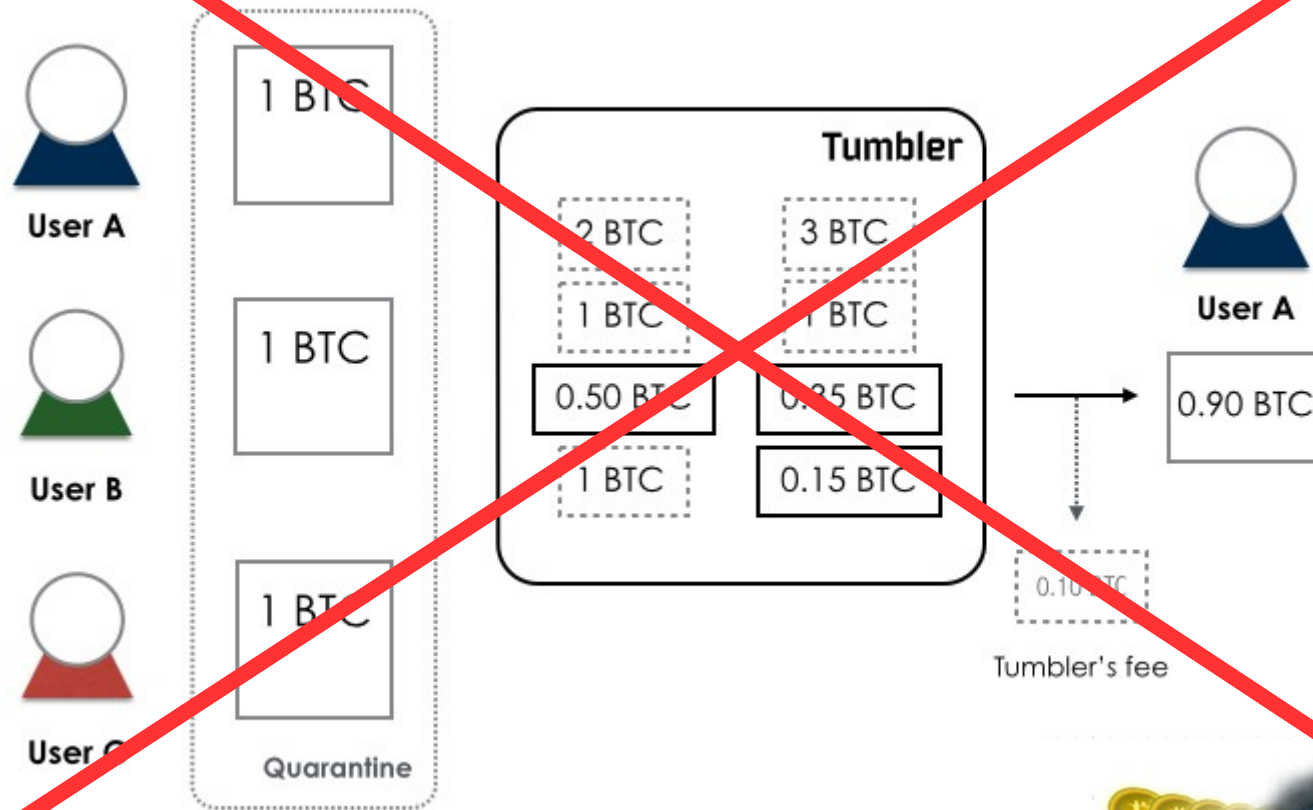
But... how?

- Mixing Services



But... how?

- ~~Mixing Services~~



ZeroCoin

or “how to prove you have money without showing it”



Alice



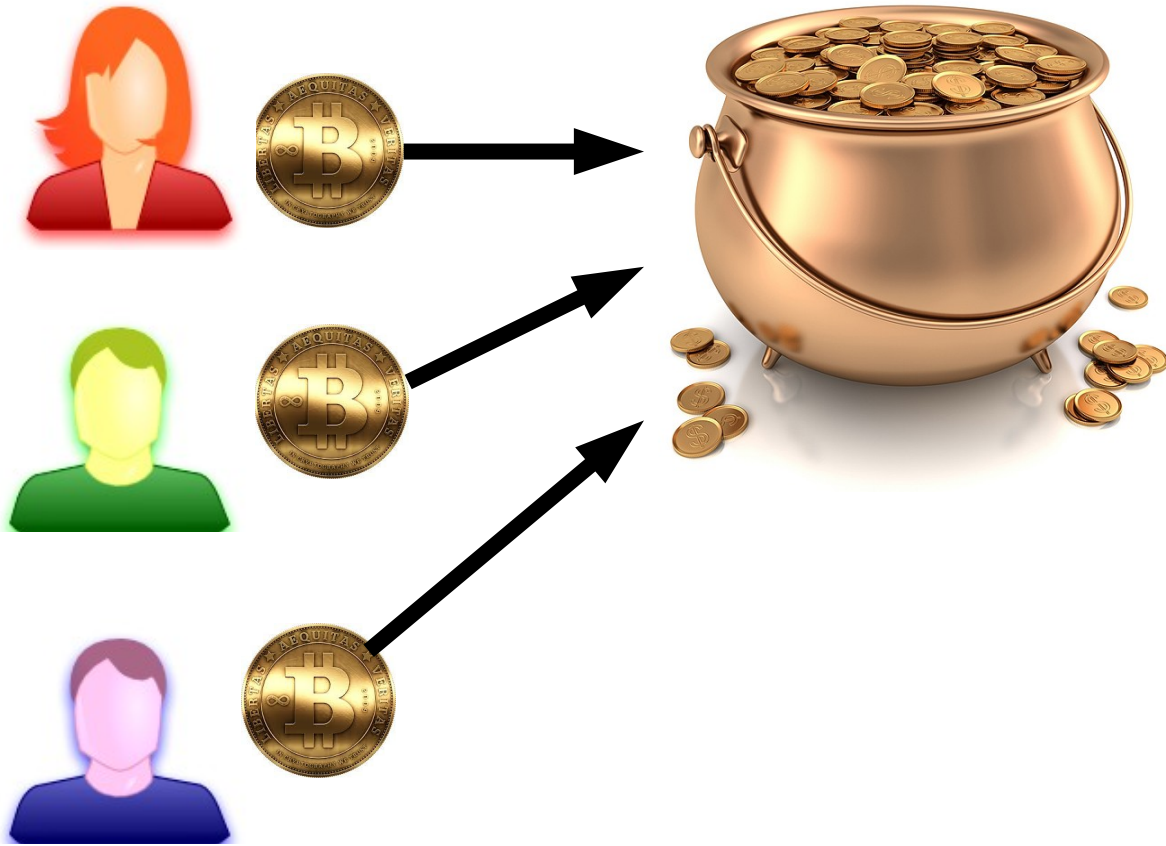
Bob



Alice

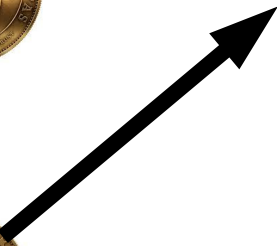
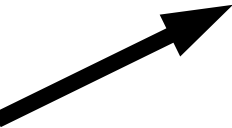
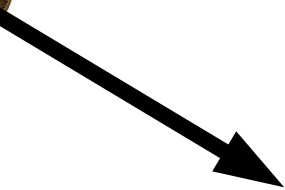


Bob



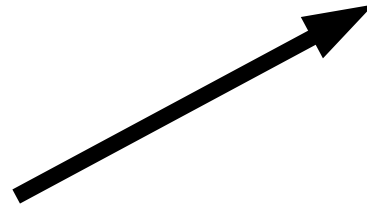
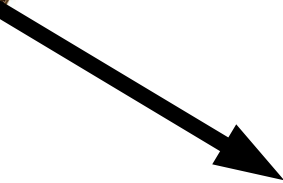


Alice





Alice



Bob



Alice



Bob



Basic Idea

- 4 operations
 - setup()
 - **mint()**
 - **spend()**
 - verify()

mint()



mint()



1 BTC



mint()



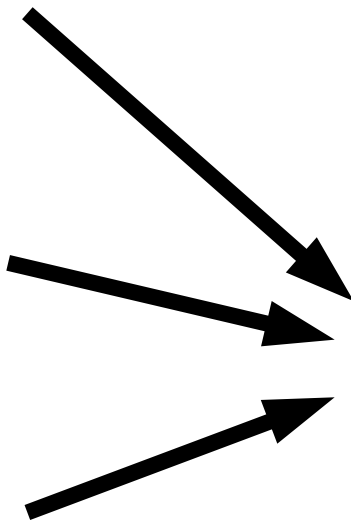
1 BTC



mint()



1 BTC



mint()



mint()



1 BTC



mint()



1 BTC



mint()



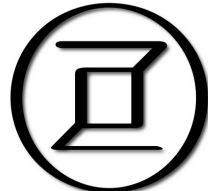
1 BTC



Alice



mint()



1 BTC

1 ZC



mint()



mint()



1 BTC



mint()



1 BTC



mint()



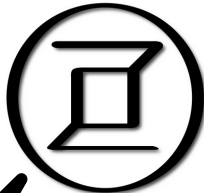
1 BTC



Alice



mint()



1 BTC

1 ZC

pub_{zc}



spend()



Alice

spend()

proof

$$(f+g)'(x) = f'(x) + g'(x)$$

proof from book:

$$(f+g)'(x) = \lim_{h \rightarrow 0} \frac{(f+g)(x+h) - (f+g)(x)}{h}$$
$$\rightarrow = \lim_{h \rightarrow 0} \frac{[f(x+h) + g(x+h)] - [f(x) + g(x)]}{h}$$
$$\rightarrow = \lim_{h \rightarrow 0} \frac{[f(x+h) - f(x)] + [g(x+h) - g(x)]}{h}$$
$$\rightarrow = \lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h} + \lim_{h \rightarrow 0} \frac{g(x+h) - g(x)}{h}$$

how does this happen?

$$\rightarrow = f'(x) + g'(x)$$

why isn't

$$\lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h}$$


zero? $f(x) - f(x)$?



Alice

spend()

proof



Alice

$$(f+g)'(x) = f'(x) + g'(x)$$

proof from book:

$$(f+g)'(x) = \lim_{h \rightarrow 0} \frac{(f+g)(x+h) - (f+g)(x)}{h}$$
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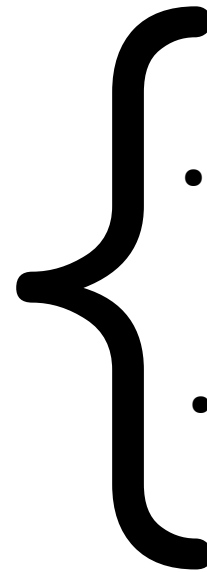
how does this happen

$$\rightarrow = \lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h} + \lim_{h \rightarrow 0} \frac{g(x+h) - g(x)}{h}$$
$$\rightarrow = f'(x) + g'(x)$$

why isn't

$$\lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h}$$


zeros? $f(x) - f(x)$?



- “I've minted one ZC”
- “I haven't spend it”

spend()

proof



Alice

$$(f+g)'(x) = f'(x) + g'(x)$$

proof from book:

$$(f+g)'(x) = \lim_{h \rightarrow 0} \frac{(f+g)(x+h) - (f+g)(x)}{h}$$
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how does this happen

$$\rightarrow = \lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h} + \lim_{h \rightarrow 0} \frac{g(x+h) - g(x)}{h}$$
$$\rightarrow = f'(x) + g'(x)$$

why isn't

$$\lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h}$$

zeros? $f(x) - f(x)$?

- “I know one (unspent) pub_{zc} ”
- “I know the construction of pub_{zc} ”

spend()



$$(f+g)'(x) = f'(x) + g'(x)$$

proof from book:

$$(f+g)'(x) = \lim_{h \rightarrow 0} \frac{(f+g)(x+h) - (f+g)(x)}{h}$$
$$\rightarrow \lim_{h \rightarrow 0} \frac{[f(x+h) + g(x+h)] - [f(x) + g(x)]}{h}$$
$$\rightarrow \lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h} + \lim_{h \rightarrow 0} \frac{g(x+h) - g(x)}{h}$$

how does this happen?

$$\rightarrow f'(x) + g'(x)$$

why isn't




$$\lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h} = f(x) - f(x) ?$$



Bob





 Bitcoin
  Zerocoin Mint
  Zerocoin Spend

verify()



Challenges?

Cryptographic Building Blocks

- Commitment Scheme
- Zero-Knowledge Proofs
- Accumulator

Commitment Scheme

- “How would you flip a coin over the phone?”

1. flips coin = x



Alice



Bob

1. flips coin = x
2. random r_A



Alice



Bob

1. flips coin = x
2. random r_A
3. $h(x, r_A)$



Alice



Bob

1. flips coin = x

2. random r_A

3. $h(x, r_A)$



Alice



Bob

4. chooses y



1. flips coin = x

2. random r_A

3. $h(x, r_A)$



4. chooses y

A horizontal black arrow pointing from the right towards the left, indicating the direction of communication from Bob to Alice.

5. x, r_A

A horizontal black arrow pointing from the left towards the right, indicating the direction of communication from Alice to Bob.

1. flips coin = x

2. random r_A

3. $h(x, r_A)$



4. chooses y

A horizontal black arrow pointing from right to left, indicating the direction of communication from Bob to Alice.

5. x, r_A

A horizontal black arrow pointing from left to right, indicating the direction of communication from Alice to Bob.

$$h(x, r_A) == h(y, r_A)$$



Pedersen Commitment Scheme (*)

- group $G = \langle g \rangle = \langle h \rangle$

Pedersen Commitment Scheme (*)

- group $G = \langle g \rangle = \langle h \rangle$
- commit to value s :

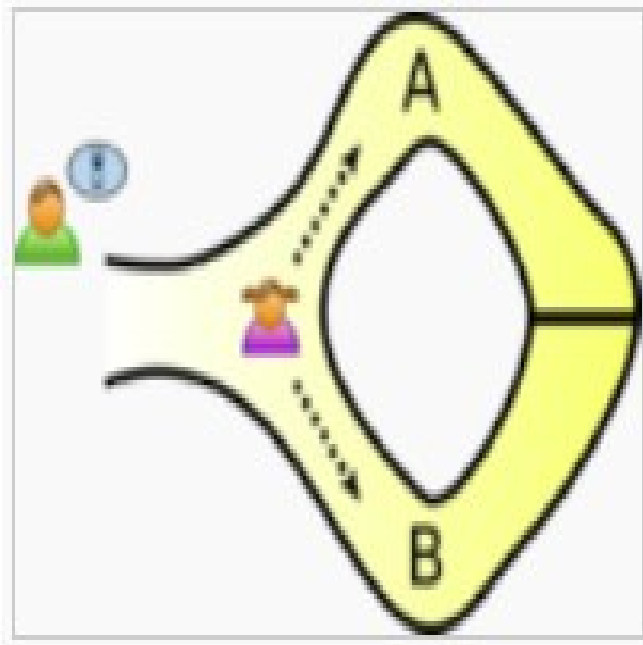
choose random r

$$\text{pub} = g^s h^r$$

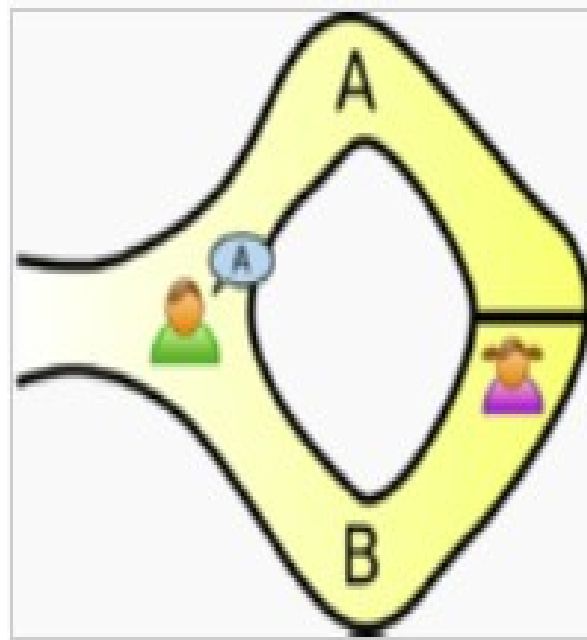
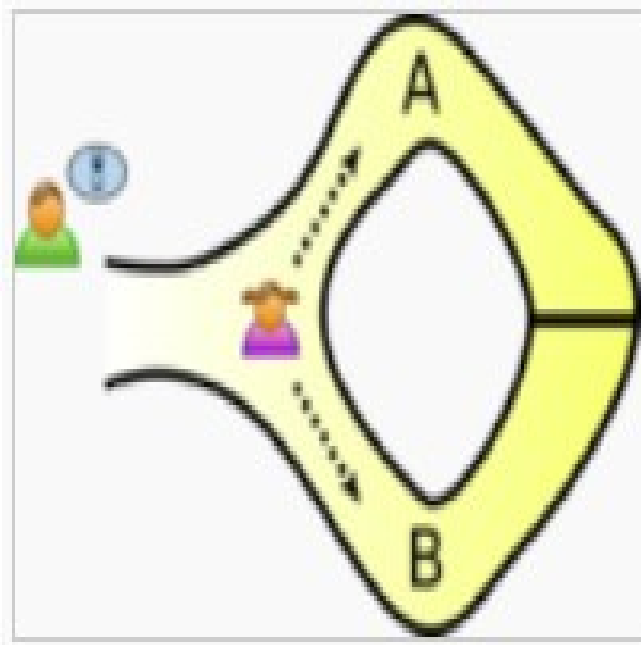
$$\text{sec} = (s, r)$$

Zero-Knowledge Proofs

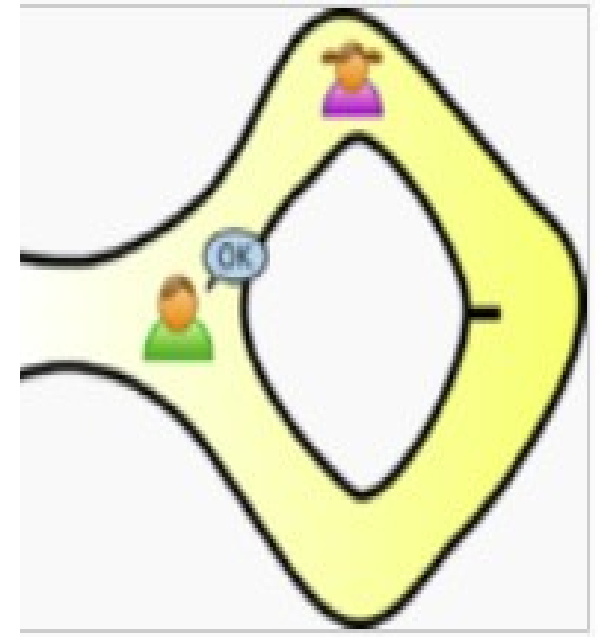
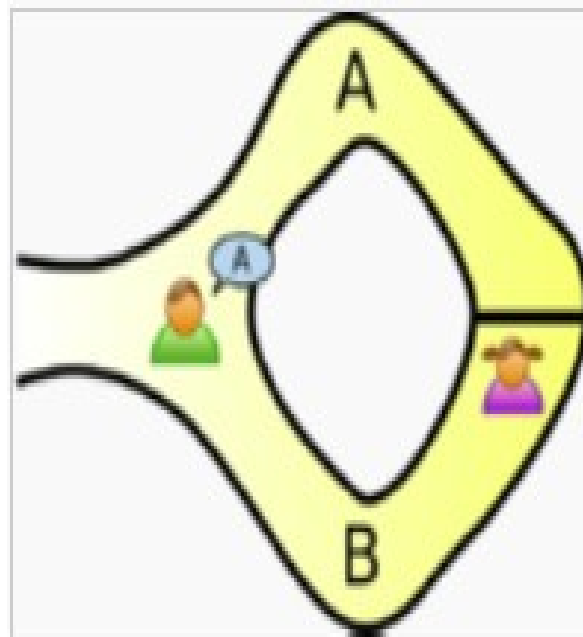
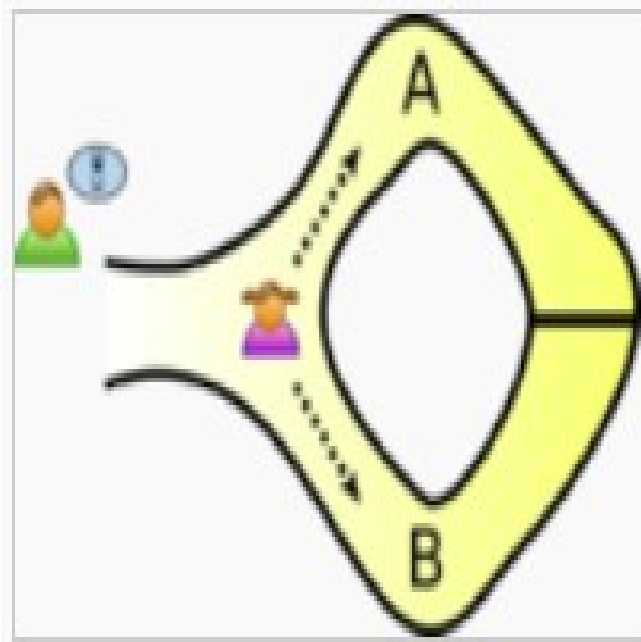
Interactive Zero-Knowledge Proofs



Interactive Zero-Knowledge Proofs



Interactive Zero-Knowledge Proofs



Non-Interactive Zero-Knowledge Proofs

- Fiat-Shamir Heuristic
- can be used as a Signature of Knowledge (**ZKSoK**)

Accumulator

- **Given:** $\mathbf{C} := \{ \text{pub}_{zc, i} \mid i = 1, \dots, n \}$
- **Show:** “I know $\text{pub}_{zc, j} \in \mathbf{C}$ ” without revealing it

Accumulator

- **Given:** $\mathbf{C} := \{ \text{pub}_{zc, i} \mid i = 1, \dots, n \}$
- **Show:** “I know $\text{pub}_{zc, j} \in \mathbf{C}$ ” without revealing it

$$(\text{pub}_{zc, j} = \text{pub}_{zc, 1}) \vee \dots \vee (\text{pub}_{zc, j} = \text{pub}_{zc, n})$$

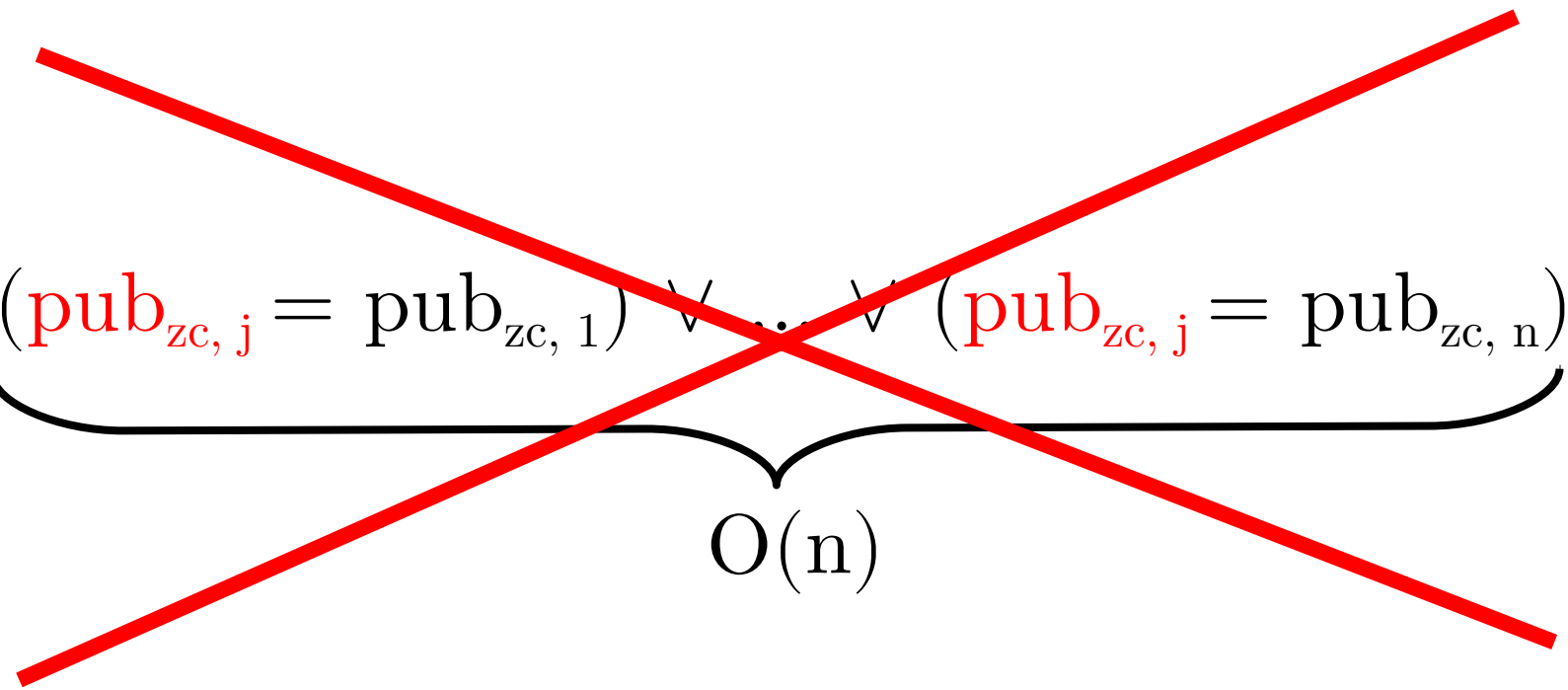
Accumulator

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$$\underbrace{(\text{pub}_{zc, j} = \text{pub}_{zc, 1}) \vee \dots \vee (\text{pub}_{zc, j} = \text{pub}_{zc, n})}_{O(n)}$$

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public one-way accumulator:

$$\text{Acc} = \text{accumulate}(\mathbf{C})$$

Accumulator

- **Given:** $\mathbf{C} := \{ \text{pub}_{zc, i} \mid i = 1, \dots, n \}$
- **Show:** “I know $\text{pub}_{zc, j} \in \mathbf{C}$ ” without revealing it

public one-way accumulator:

$$\text{Acc} = \text{accumulate}(\mathbf{C})$$

$$\text{wit} = \text{generateWitness}(\mathbf{C}, \text{value})$$

Accumulator

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public one-way accumulator:

$\text{Acc} = \text{accumulate}(\mathbf{C})$

$\text{wit} = \text{generateWitness}(\mathbf{C}, \text{value})$

$\text{accVerify}(\text{Acc}, \text{value}, \text{wit}) \rightarrow \{0, 1\}$

Accumulator

- **Given:** $\mathbf{C} := \{ \text{pub}_{zc, i} \mid i = 1, \dots, n \}$
- **Show:** “I know $\text{pub}_{zc, j} \in \mathbf{C}$ ” without revealing it

public one-way accumulator:

$$\text{Acc} = \text{accumulate}(\mathbf{C}) = \prod_{i=1}^n \text{pub}_{zc, i}$$

Accumulator

- **Given:** $\mathbf{C} := \{ \text{pub}_{zc, i} \mid i = 1, \dots, n \}$
- **Show:** “I know $\text{pub}_{zc, j} \in \mathbf{C}$ ” **without revealing it**

public one-way accumulator:

$$\text{Acc} = \text{accumulate}(\mathbf{C}) = \prod_{i=1}^n \text{pub}_{zc, i}$$

Accumulator (*)

- **Given:** $\mathbf{C} := \{ \text{pub}_{zc, i} \mid i = 1, \dots, n \}$
- **Show:** “I know $\text{pub}_{zc, j} \in \mathbf{C}$ ” **without revealing it**

public one-way accumulator:

$$\text{Acc} = \text{accumulate}(\mathbf{C}) = \mathbf{u} \prod_{i=1}^n \text{pub}_{zc, i}$$

$$\text{Acc} = \text{accumulate}(\mathbf{C}) = \mathbf{u} \prod_{i=1}^n \text{pub}_{z_c, i}$$

$$\text{Acc} = \text{accumulate}(\mathbf{C}) = \mathbf{u} \prod_{i=1}^n \text{pub}_{z_c, i}$$

$$\text{wit} = \text{generateWitness}(\mathbf{C}, \text{value}) = \text{accumulate}(\mathbf{C} \setminus \{\text{value}\})$$

$$\text{Acc} = \text{accumulate}(\mathbf{C}) = \mathbf{u} \prod_{i=1}^n \text{pub}_{z_c, i}$$

$$\text{wit} = \text{generateWitness}(\mathbf{C}, \text{value}) = \text{accumulate}(\mathbf{C} \setminus \{\text{value}\})$$

$$\text{accVerify}(\text{Acc}, \text{value}, \text{wit}) = 1 \text{ iff } \text{wit}^{\text{value}} = \text{Acc}$$

ZeroCoin Protocol

- 4 operations
 - setup()
 - **mint()**
 - **spend()**
 - verify()

setup()

- setup accumulator
- setup commitment parameters

$$G = \langle g \rangle = \langle h \rangle$$

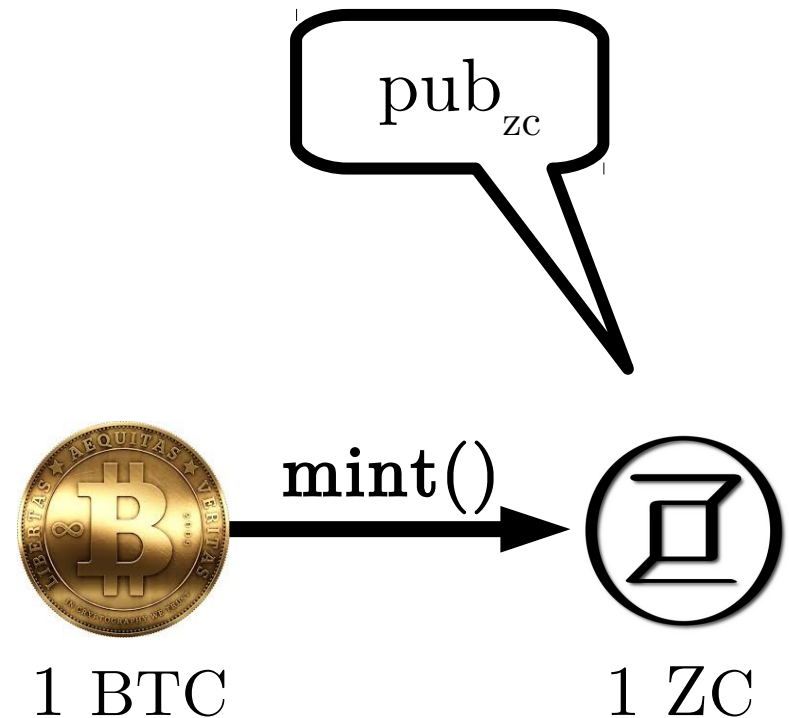
$$\text{mint}(I_{\text{btc}}) \rightarrow (\text{pub}_{\text{zc}}, \text{sec}_{\text{zc}})$$

- commit to value s :

choose random r

$$\text{pub}_{\text{zc}} = g^s h^r$$

$$\text{sec}_{\text{zc}} = (s, r)$$



$$\text{spend}(\text{pub}_{zc}, \text{sec}_{zc}, \mathbf{C}, \mathbf{O}_{\text{btc}}) \rightarrow (\pi, \mathbf{s})$$

- $\text{Acc} = \text{accumulate}(\mathbf{C})$

$$\text{spend}(\text{pub}_{zc}, \text{sec}_{zc}, \mathbf{C}, \mathbf{O}_{\text{btc}}) \rightarrow (\pi, s)$$

- $\text{Acc} = \text{accumulate}(\mathbf{C})$
- $\text{wit} = \text{generateWitness}(\mathbf{C}, \text{pub}_{zc})$

$$\text{spend}(\text{pub}_{zc}, \text{sec}_{zc}, \mathbf{C}, \mathbf{O}_{\text{btc}}) \rightarrow (\pi, \mathbf{s})$$

- $\text{Acc} = \text{accumulate}(\mathbf{C})$
- $\text{wit} = \text{generateWitness}(\mathbf{C}, \text{pub}_{zc})$
- \mathbf{s} in $\text{sec}_{zc} = (\mathbf{s}, \mathbf{r})$

$$\text{spend}(\text{pub}_{zc}, \text{sec}_{zc}, \mathbf{C}, \mathbf{O}_{\text{btc}}) \rightarrow (\pi, s)$$

- $\text{Acc} = \text{accumulate}(\mathbf{C})$
- $\text{wit} = \text{generateWitness}(\mathbf{C}, \text{pub}_{zc})$
- $s \text{ in } \text{sec}_{zc} = (s, r)$
- $\pi = \text{ZKSoK} [\mathbf{O}_{\text{btc}}] \{\dots\}$

$$\text{spend}(\text{pub}_{zc}, \text{sec}_{zc}, \mathbf{C}, \mathbf{O}_{\text{btc}}) \rightarrow (\pi, s)$$

- $\text{Acc} = \text{accumulate}(\mathbf{C})$
- $\text{wit} = \text{generateWitness}(\mathbf{C}, \text{pub}_{zc})$
- $s \text{ in } \text{sec}_{zc} = (s, r)$
- $\pi = \text{ZKSoK} [\mathbf{O}_{\text{btc}}] \{ (\text{pub}_{zc}, \text{wit}, r) :$

$$\underbrace{\text{AccVerify}(\text{Acc}, \text{pub}_{zc}, \text{wit}) = 1 \wedge \dots}$$

“I know one pub_{zc} in \mathbf{C} ”

$$\text{spend}(\text{pub}_{zc}, \text{sec}_{zc}, \mathbf{C}, \mathbf{O}_{\text{btc}}) \rightarrow (\pi, s)$$

- $\text{Acc} = \text{accumulate}(\mathbf{C})$
- $\text{wit} = \text{generateWitness}(\mathbf{C}, \text{pub}_{zc})$
- $s \text{ in } \text{sec}_{zc} = (s, r)$
- $\pi = \text{ZKSoK} [\mathbf{O}_{\text{btc}}] \{ (\text{pub}_{zc}, \text{wit}, r) :$

$$\underbrace{\text{AccVerify}(\text{Acc}, \text{pub}_{zc}, \text{wit}) = 1}_{\text{“I know one pub}_{zc} \text{ in } \mathbf{C}”} \wedge \underbrace{\text{pub}_{zc} = g^s h^r}_{\text{“I know its construction”}}$$

“I know one pub_{zc} in \mathbf{C} ”

“I know its construction”

$$\text{verify}(\pi, s, O_{\text{btc}}, \mathbf{C}) \rightarrow \{0, 1\}$$

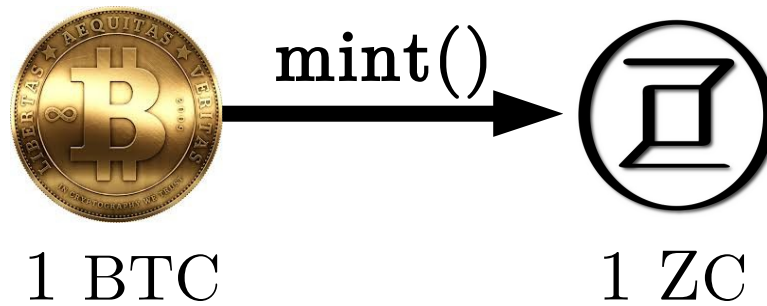
- verify if s unspent
- verify correctness of π as a signature on O_{btc}

Remarks

- accumulator checkpoint
- proof size \rightarrow memory issues
- proof complexity \rightarrow longer verification time

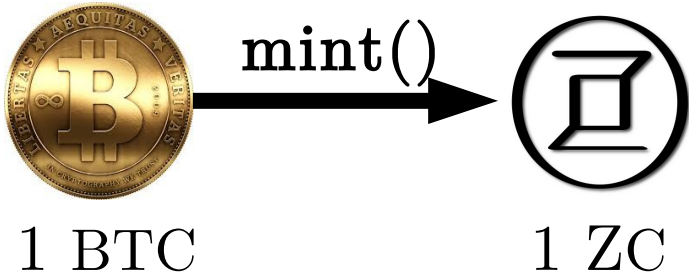
One more thing...

Is ZeroCoin enough?

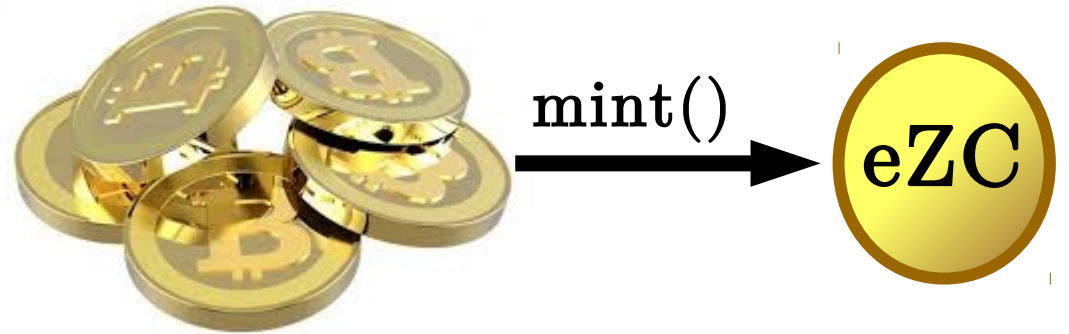
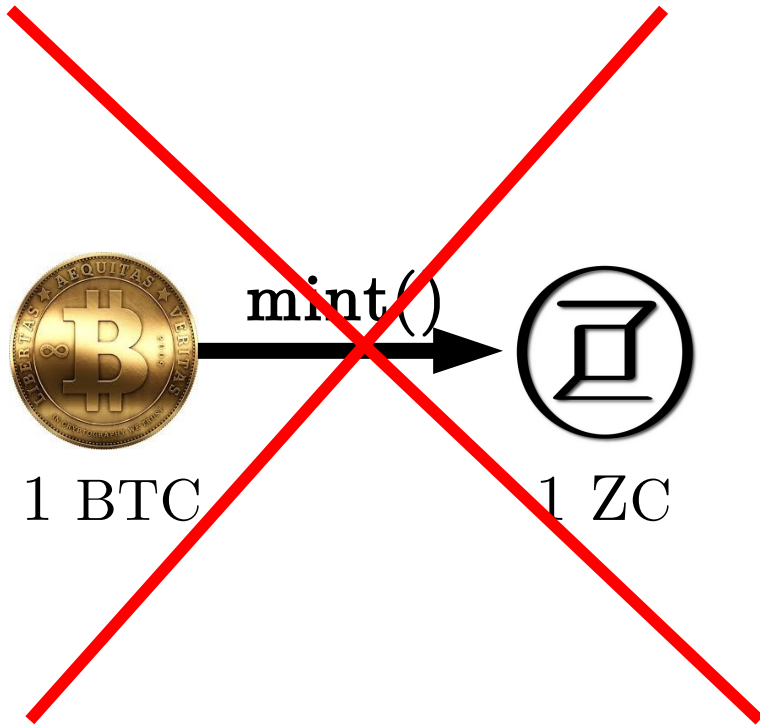


eZC: ZeroCoin Reloaded

New Ideas



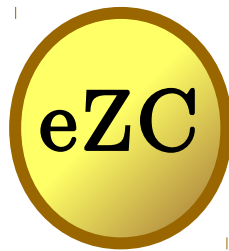
New Ideas



New Ideas



Alice



amount?

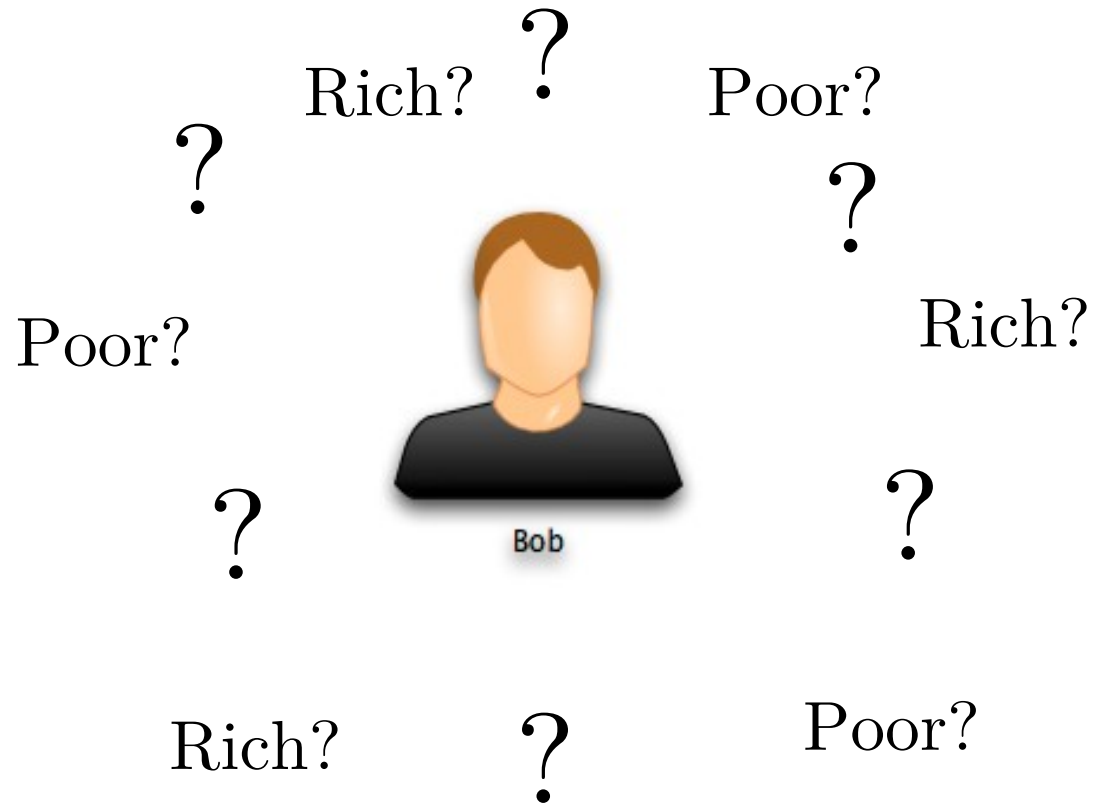


Bob

amount = 10 BTC

amount = 10 BTC

New Ideas



But... How?

- 5 operations:
 - setup()
 - **mint(amount)**
 - spendEZCtoBTC()
 - **spendEZCtoEZC()**
 - verify()

setup()

$$G = \langle g \rangle = \langle h \rangle = \langle w \rangle$$

$$\text{mint}(\mathbf{I}_{\text{btc}}) \rightarrow (\pi_{\text{pub}}, \text{pub}_{\text{ezc}}, \text{sec}_{\text{ezc}})$$

- commit to value s and transaction amount (a):

$$\text{pub}_{\text{ezc}} = g^s h^r w^a$$

$$\text{sec}_{\text{ezc}} = (s, r)$$

$$\pi_{\text{pub}} = \text{ZKPoK}\{(s, r) : \text{pub}_{\text{ezc}} = g^s h^r w^a\}$$

spendEZCtoEZC



Alice

1. Proves eZC_{send} 's validity

2. Mints eZC_{change}

3. Pre-mints eZC_{receive}



Bob

4. Mints eZC_{receive}

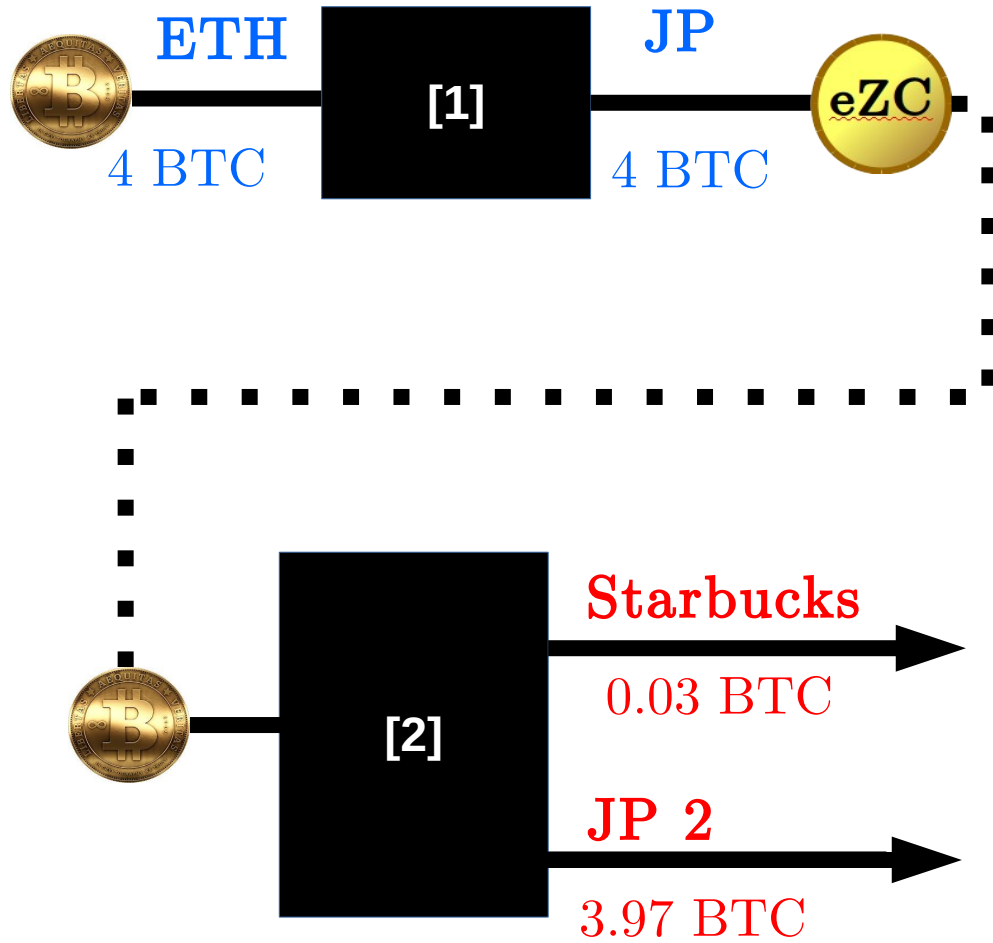
Transactions Revisited



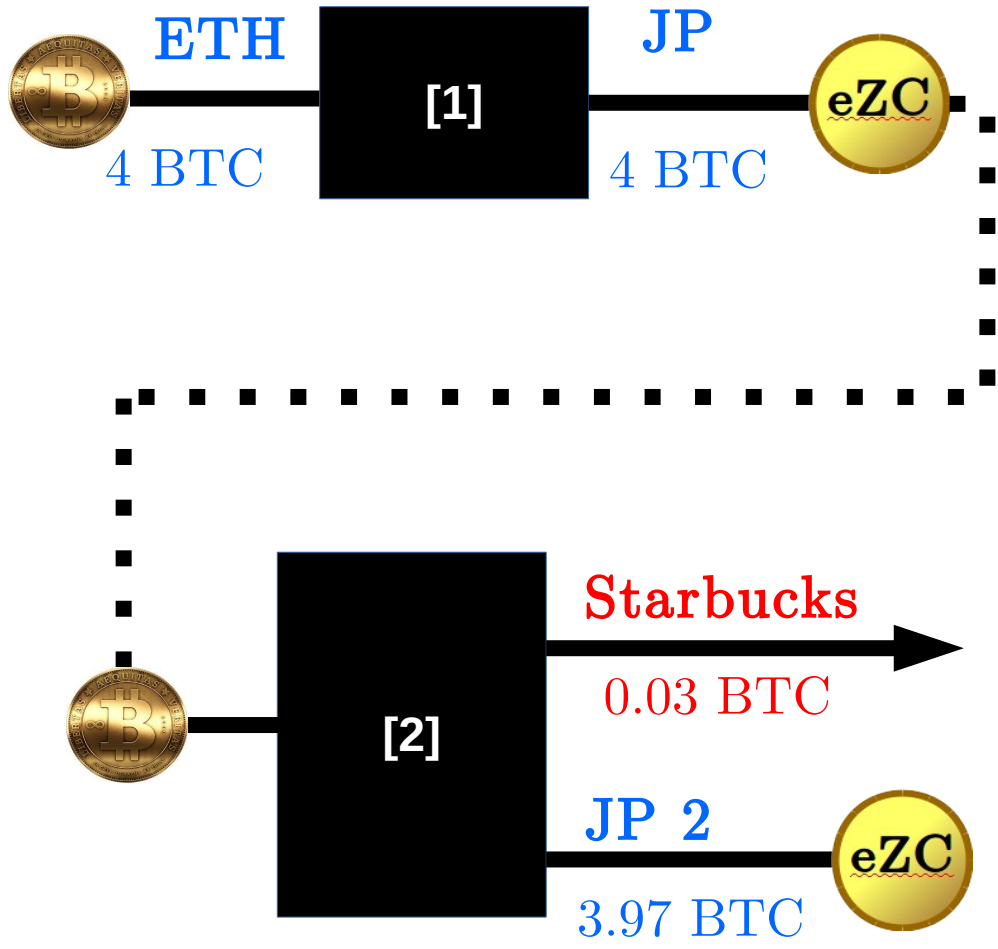
Transactions Revisited



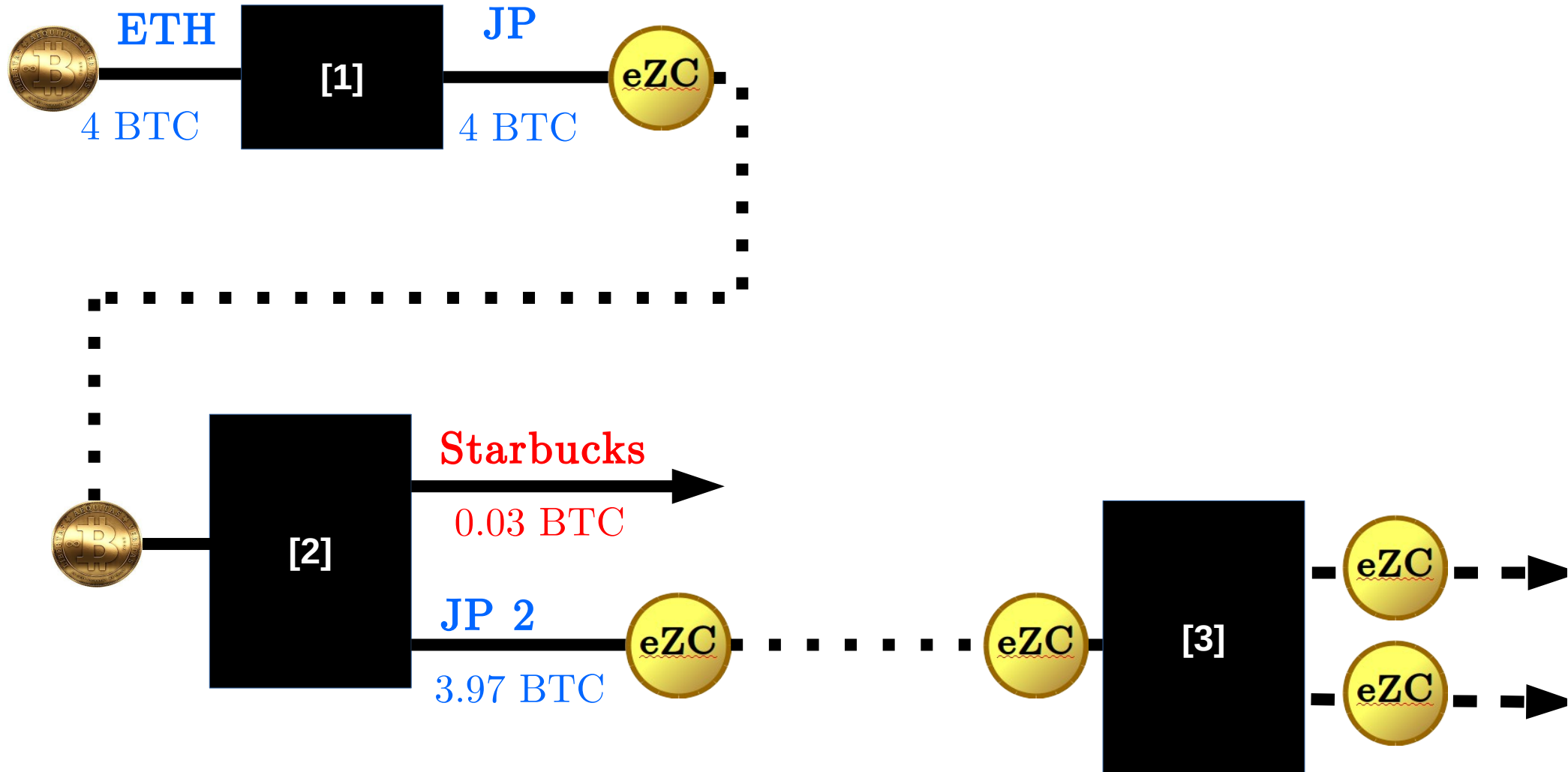
Transactions Revisited



Transactions Revisited



Transactions Revisited



Conclusion

Thank You



Tip if you enjoyed it!