Clock Synchronization

Chapter 9

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Clock Synchronization

Rating

- Area maturity

First steps

Text book

- Practical importance

No apps

Mission critical

- Theory appeal

Booooooooring

Exciting

Overview

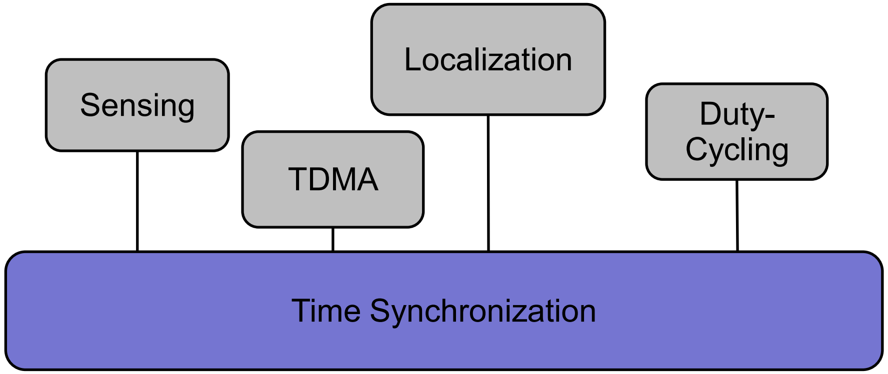
- Motivation
- Clock Sources & Hardware
- Single-Hop Clock Synchronization
- Clock Synchronization in Networks
- Protocols: RBS, TPSN, FTSP, GTSP
- Theory of Clock Synchronization
- Protocol: PulseSync

Motivation

- Synchronizing time is essential for many applications
  - Coordination of wake-up and sleeping times (energy efficiency)
  - TDMA schedules
  - Ordering of collected sensor data/events
  - Co-operation of multiple sensor nodes
  - Estimation of position information (e.g. shooter detection)

- Goals of clock synchronization
  - Compensate *offset** between clocks
  - Compensate *drift** between clocks

*terms are explained on following slides



Properties of Clock Synchronization Algorithms

- External versus internal synchronization
  - External sync: Nodes synchronize with an external clock source (UTC)
  - Internal sync: Nodes synchronize to a common time
    - to a leader, to an averaged time, or to anything else
- One-shot versus continuous synchronization
  - Periodic synchronization required to compensate clock drift
- A-priori versus a-posteriori
  - A-posteriori clock synchronization triggered by an event
- Global versus local synchronization (explained later)
- Accuracy versus convergence time, Byzantine nodes, …

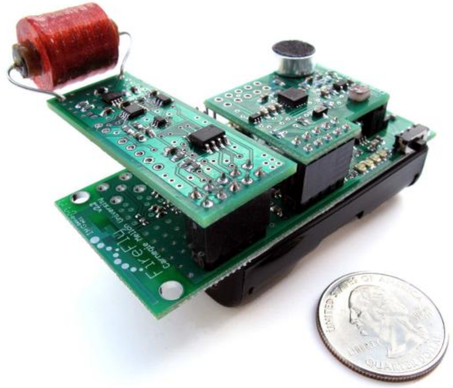
Clock Sources

- Radio Clock Signal:
  - Clock signal from a reference source (atomic clock) is transmitted over a long wave radio signal
  - DCF77 station near Frankfurt, Germany transmits at 77.5 kHz with a transmission range of up to 2000 km
  - Accuracy limited by the distance to the sender, Frankfurt-Zurich is about 1ms.
  - Special antenna/receiver hardware required

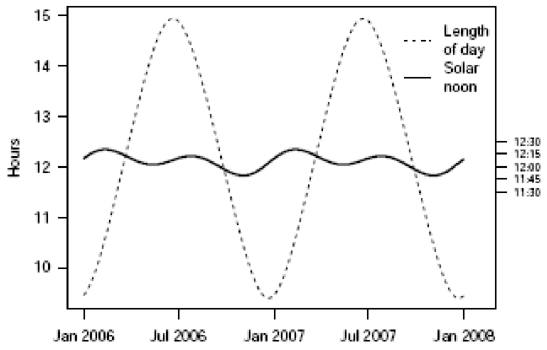
- Global Positioning System (GPS):
  - Satellites continuously transmit own position and time code
  - Line of sight between satellite and receiver required
  - Special antenna/receiver hardware required

Clock Sources (2)

- AC power lines:
  - Use the magnetic field radiating from electric AC power lines
  - AC power line oscillations are extremely stable (10^{-8} ppm)
  - Power efficient, consumes only 58 µW
  - Single communication round required to correct phase offset after initialization

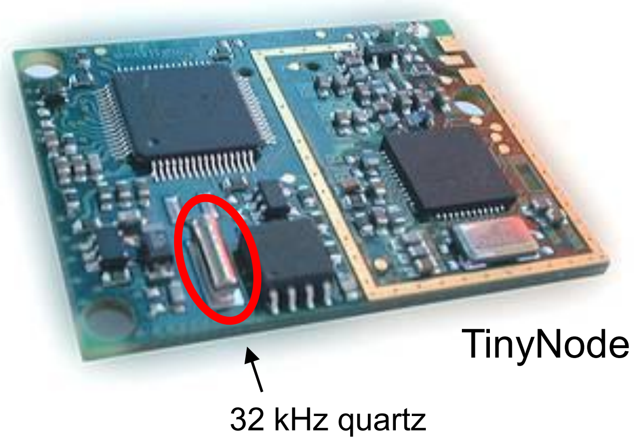


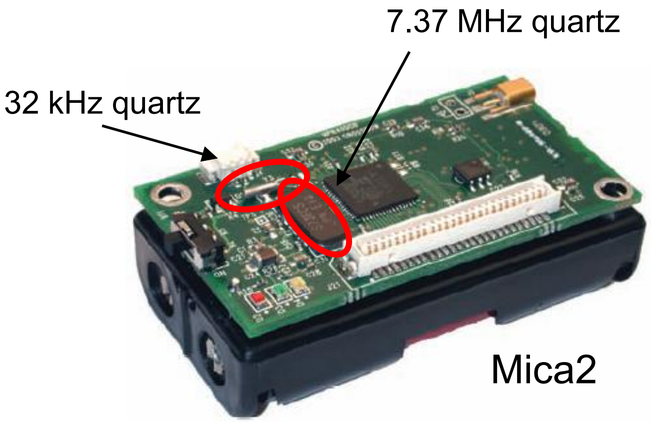
- Sunlight:
  - Using a light sensor to measure the length of a day
  - Offline algorithm for reconstructing global timestamps by correlating annual solar patterns (no communication required)



Clock Devices in Sensor Nodes

- Structure
  - External oscillator with a nominal frequency (e.g. 32 kHz or 7.37 MHz)
  - Counter register which is incremented with oscillator pulses
  - Works also when CPU is in sleep state

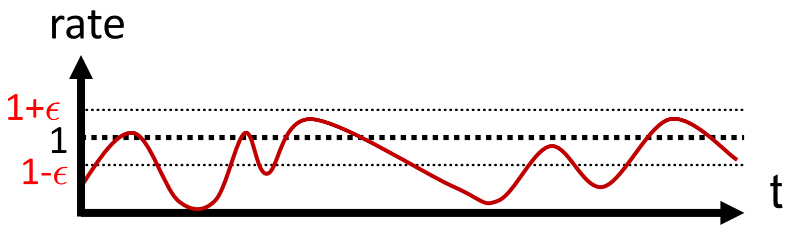




Platform	System clock	Crystal oscillator
Mica2	7.37 MHz	32 kHz, 7.37 MHz
TinyNode 584	8 MHz	32 kHz
Tmote Sky	8 MHz	32 kHz

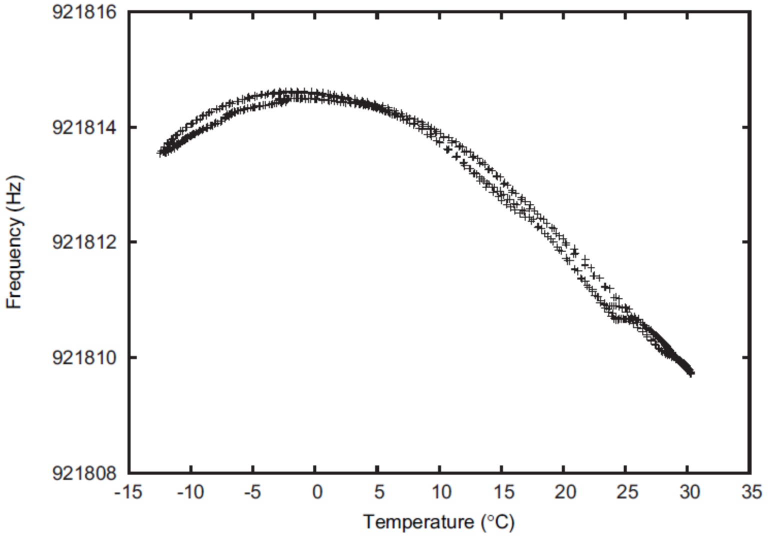
Clock Drift

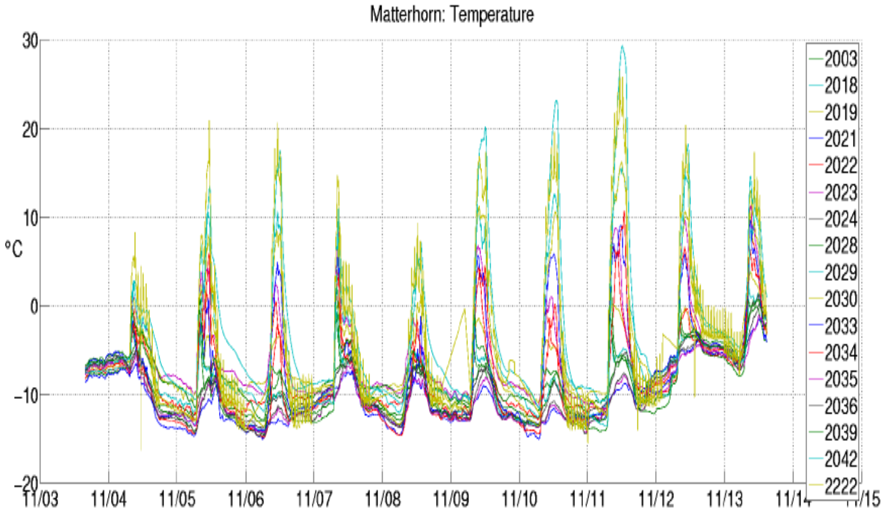
- Accuracy
  - Clock drift: random deviation from the nominal rate dependent on power supply, temperature, etc.



This is a drift of up to 50 µs per second or 0.18s per hour

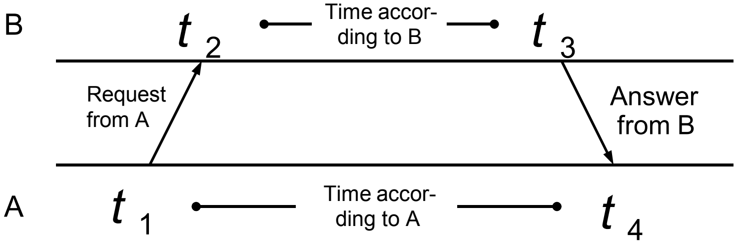
  - E.g. TinyNodes have a maximum drift of 30-50 ppm at room temperature





Sender/Receiver Synchronization

- Round-Trip Time (RTT) based synchronization



- Receiver synchronizes to the sender„s clock
- Propagation delay δ and clock offset θ can be calculated

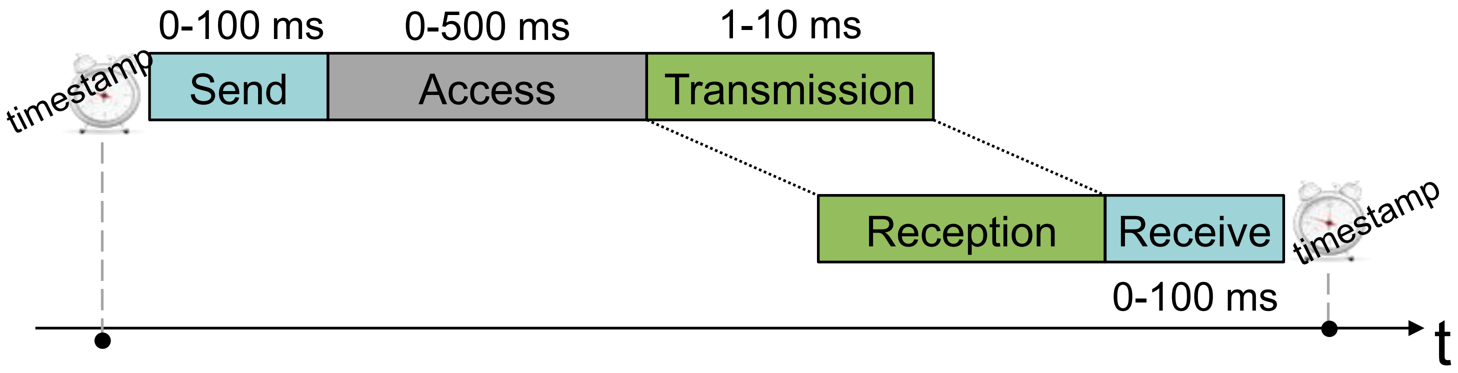
$$\delta = \frac{(t_4 - t_1) - (t_3 - t_2)}{2}$$

$$\theta = \frac{(t_2 - (t_1 + \delta)) - (t_4 - (t_3 + \delta))}{2} = \frac{(t_2 - t_1) + (t_3 - t_4)}{2}$$

Messages Experience Jitter in the Delay

- Problem: Jitter in the message delay

  Various sources of errors (deterministic and non-deterministic)

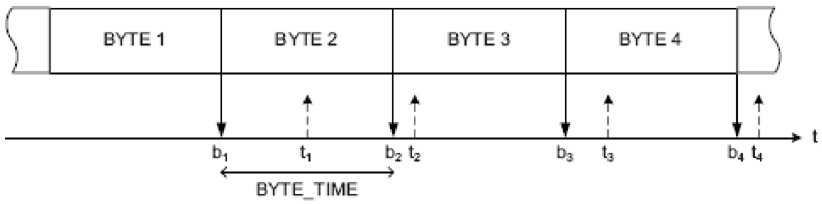


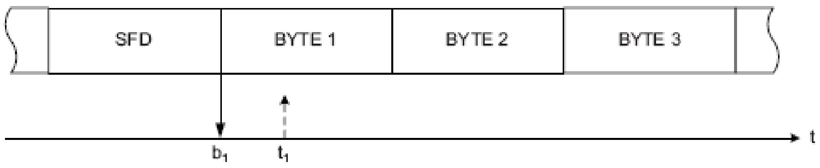
- Solution: Timestamping packets at the MAC layer (Maróti et al.)

  → Jitter in the message delay is reduced to a few clock ticks

Some Details

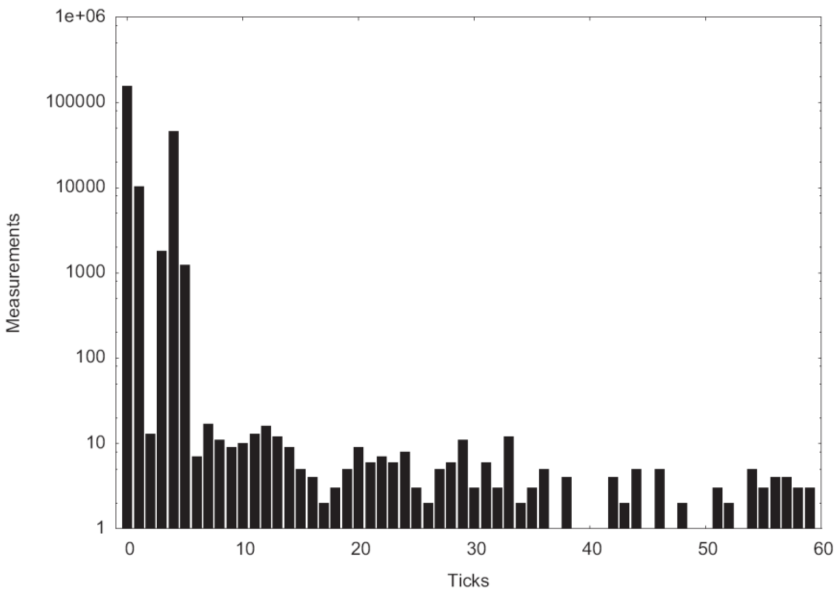
- Different radio chips use different paradigms:
  - Left is a CC1000 radio chip which generates an interrupt with each byte.
  - Right is a CC2420 radio chip that generates a single interrupt for the packet after the start frame delimiter is received.





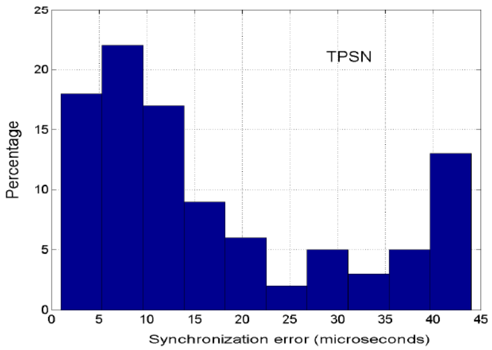
- In sensor networks propagation can be ignored (<1μs for 300m).

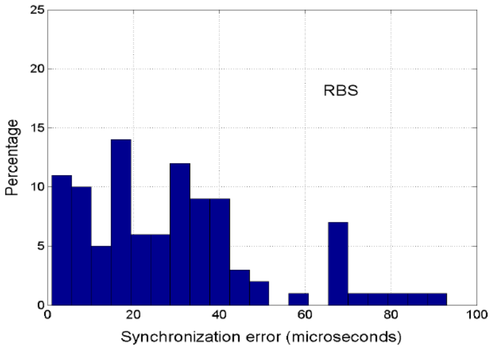
- Still there is quite some variance in transmission delay because of latencies in interrupt handling (picture right).

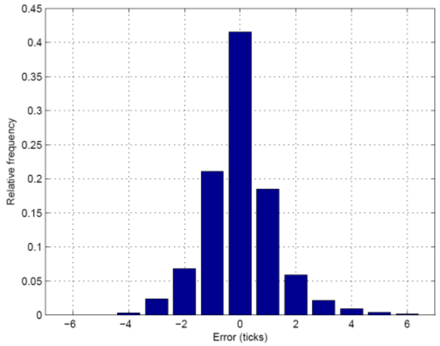


Symmetric Errors

- Many protocols don't even handle single-hop clock synchronization well. On the left figures we see the absolute synchronization errors of TPSN and RBS, respectively. The figure on the right presents a single-hop synchronization protocol minimizing systematic errors.







- Even perfectly symmetric errors will sum up over multiple hops.
  - In a chain of *n* nodes with a standard deviation σ on each hop, the expected error between head and tail of the chain is in the order of $\sigma\sqrt{n}$.

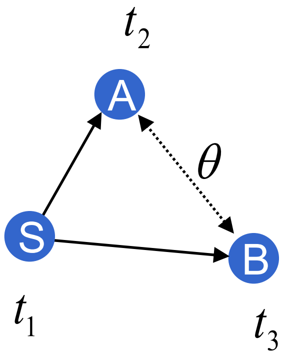
Reference-Broadcast Synchronization (RBS)

- A sender synchronizes a set of receivers with one another
- Point of reference: beacon's arrival time

$$t_2 = t_1 + S_S + A_S + P_{S,A} + R_A$$

$$t_3 = t_1 + S_S + A_S + P_{S,B} + R_B$$

$$\theta = t_2 - t_3 = (P_{S,A} - P_{S,B}) + (R_A - R_B)$$

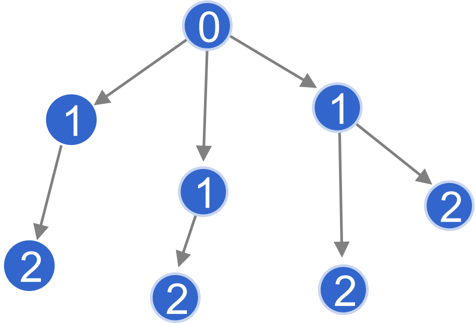


- Only sensitive to the difference in propagation and reception time
- Time stamping at the interrupt time when a beacon is received
- After a beacon is sent, all receivers exchange their reception times to calculate their clock offset

- Post-synchronization possible
- E.g., least-square linear regression to tackle clock drifts
- Multi-hop?

Time-sync Protocol for Sensor Networks (TPSN)

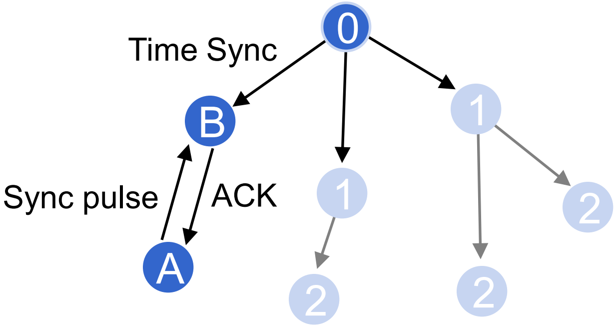
- Traditional sender-receiver synchronization (RTT-based)
- *Initialization phase: Breadth-first-search flooding*
  - Root node at level 0 sends out a *level discovery* packet
  - Receiving nodes which have not yet an assigned level set their level to +1 and start a random timer
  - After the timer is expired, a new level discovery packet will be sent
  - When a new node is deployed, it sends out a *level request* packet after a random timeout



Why this random timer?

Time-sync Protocol for Sensor Networks (TPSN)

- *Synchronization phase*
  - Root node issues a *time sync* packet which triggers a random timer at all level 1 nodes
  - After the timer is expired, the node asks its parent for synchronization using a *synchronization pulse*
  - The parent node answers with an *acknowledgement*
  - Thus, the requesting node knows the round trip time and can calculate its clock offset
  - Child nodes receiving a synchronization pulse also start a random timer themselves to trigger their own synchronization

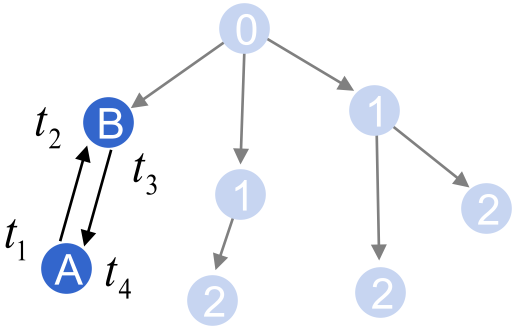


Time-sync Protocol for Sensor Networks (TPSN)

$$t_2 = t_1 + S_A + A_A + P_{A,B} + R_B$$

$$t_4 = t_3 + S_B + A_B + P_{B,A} + R_A$$

$$\theta = \frac{(S_A - S_B) + (A_A - A_B) + (P_{A,B} - P_{B,A}) + (R_B - R_A)}{2}$$

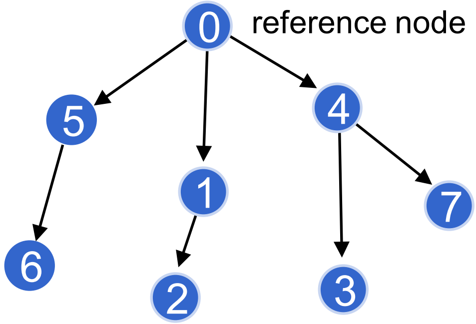


- Time stamping packets at the MAC layer
- In contrast to RBS, the signal propagation time might be negligible
- Authors claim that it is “about two times” better than RBS
- Again, clock drifts are taken into account using periodical synchronization messages

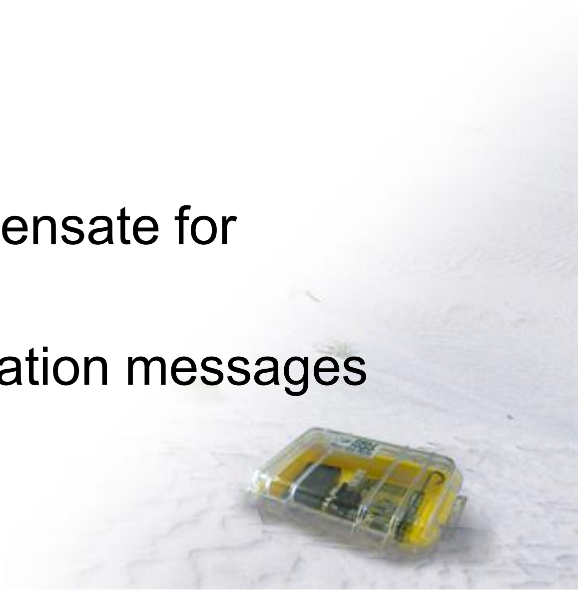
- Problem: What happens in a non-tree topology (e.g. grid)?
  - Two neighbors may have bad synchronization?

Flooding Time Synchronization Protocol (FTSP)

- Each node maintains both a local and a global time
- Global time is synchronized to the local time of a reference node
  - Node with the smallest id is elected as the reference node
- Reference time is flooded through the network periodically

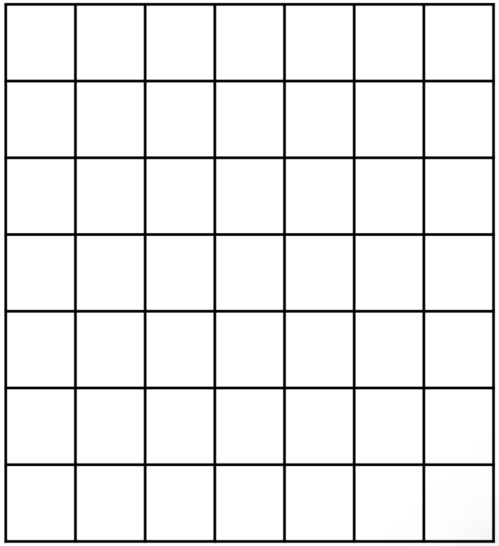


- Timestamping at the MAC Layer is used to compensate for deterministic message delays
- Compensation for clock drift between synchronization messages using a linear regression table



Best tree for tree-based clock synchronization?

- Finding a good tree for clock synchronization is a tough problem
  - Spanning tree with small (maximum or average) stretch.



- Example: Grid network, with $n = m^2$ nodes.

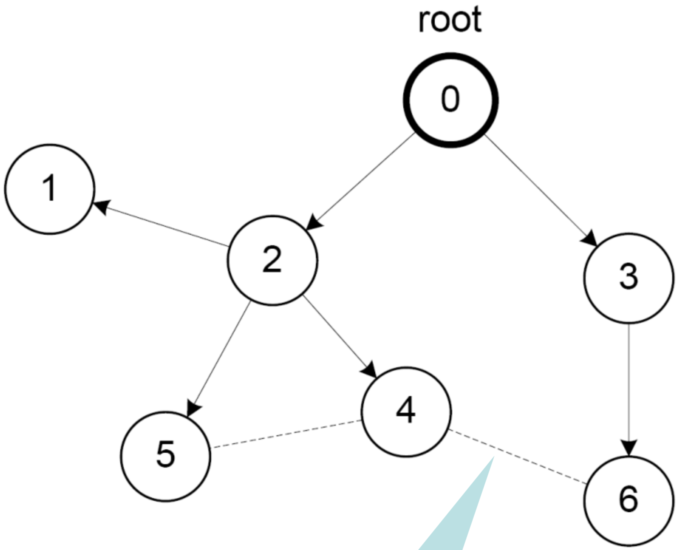
- No matter what tree you use, the maximum stretch of the spanning tree will always be at least m (just try on the grid figure right…)

- In general, finding the minimum max stretch spanning tree is a hard problem, however approximation algorithms exist [Emek, Peleg, 2004].

Variants of Clock Synchronization Algorithms

Tree-like Algorithms

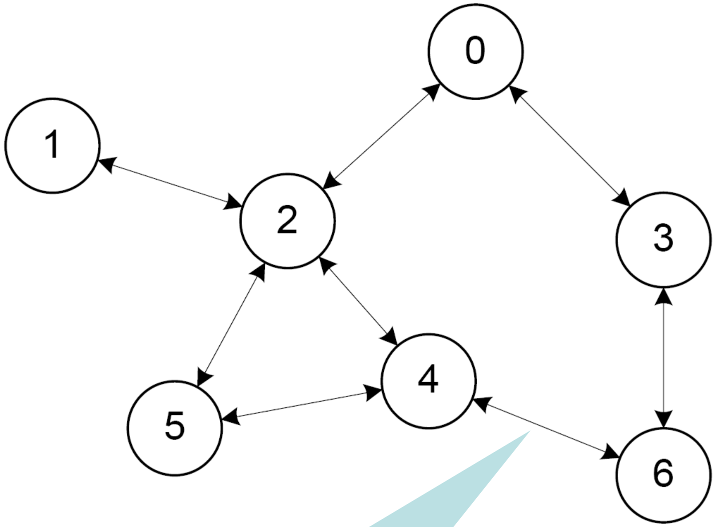
e.g. FTSP



Bad local skew

Distributed Algorithms

e.g. GTSP

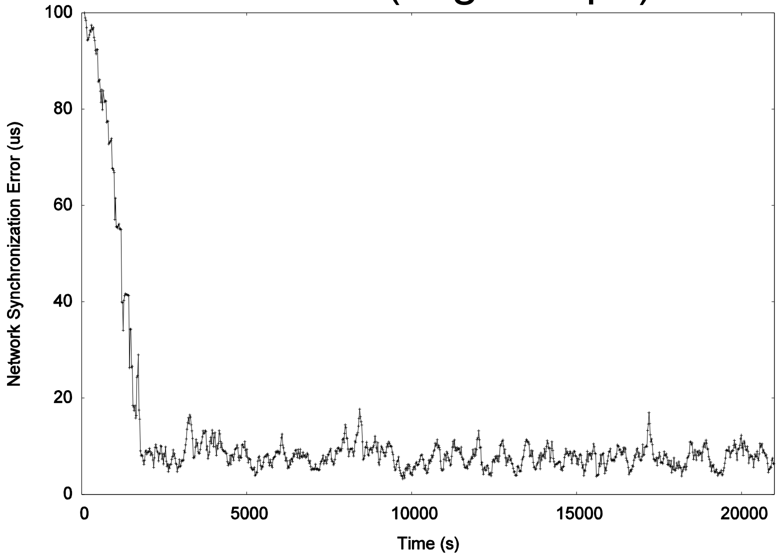


All nodes consistently average errors to *all* neigbhors

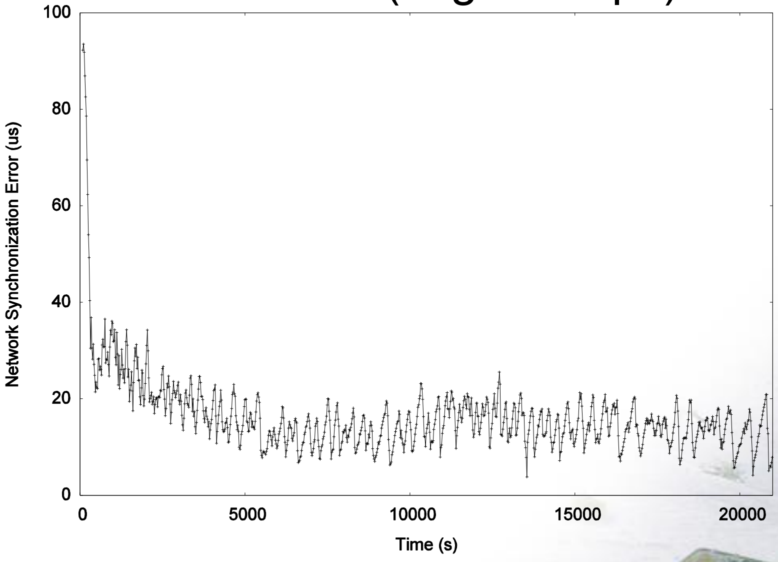
FTSP vs. GTSP: Global Skew

- Network synchronization error (global skew)
  - Pair-wise synchronization error between any two nodes in the network

FTSP (avg: 7.7 µs)



GTSP (avg: 14.0 µs)

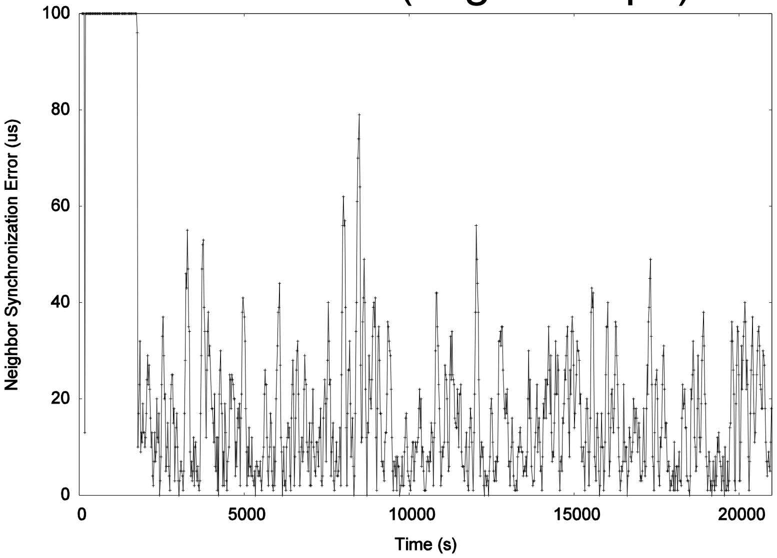


FTSP vs. GTSP: Local Skew

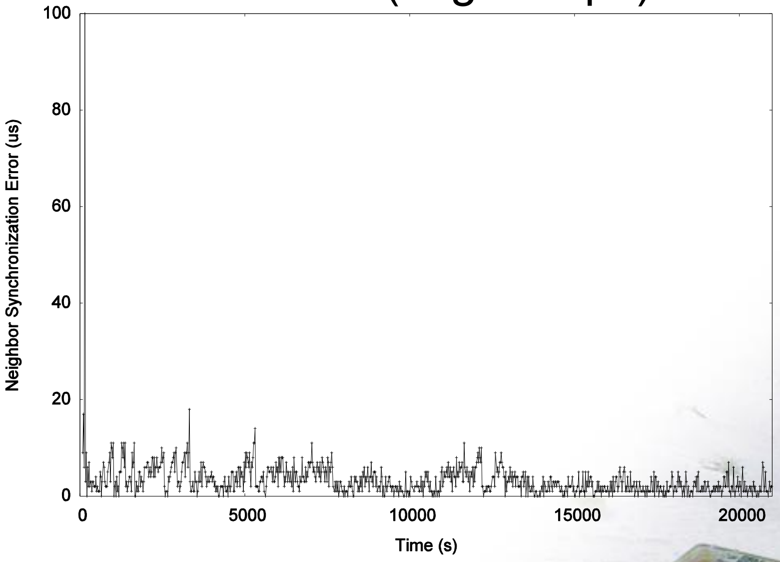
- Neighbor Synchronization error (local skew)
  - Pair-wise synchronization error between neighboring nodes

- Synchronization error between two direct neighbors:

FTSP (avg: 15.0 µs)



GTSP (avg: 2.8 µs)



Global vs. Local Time Synchronization

- Common time is essential for many applications:

Global

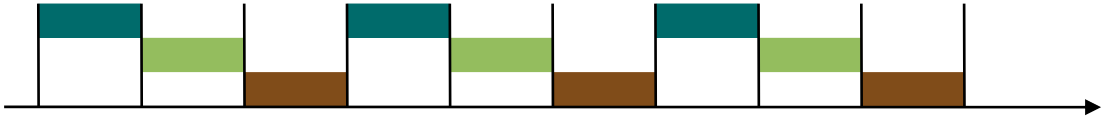
  - Assigning a timestamp to a globally sensed event (e.g. earthquake)

Local

  - Precise event localization (e.g. shooter detection, multiplayer games)

Local

  - TDMA-based MAC layer in wireless networks



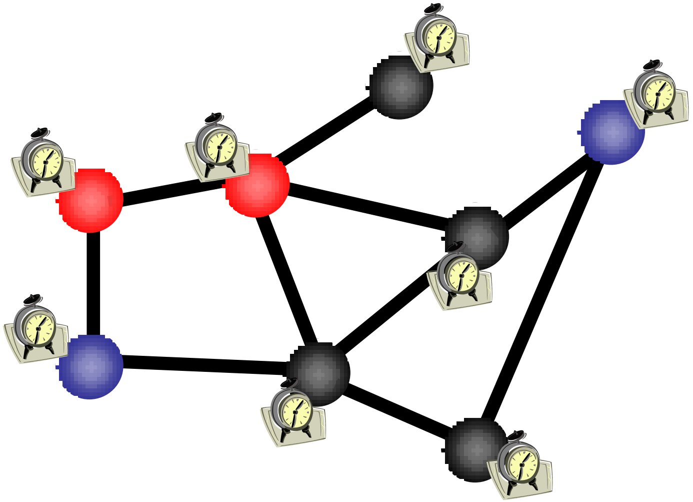
Local

  - Coordination of wake-up and sleeping times (energy efficiency)

Theory of Clock Synchronization

- Given a communication network
  1. Each node equipped with hardware clock with drift
  2. Message delays with jitter

worst-case (but constant)



- Goal: Synchronize Clocks ("Logical Clocks")
  - Both global and local synchronization!

Time Must Behave!

- Time (logical clocks) should not be allowed to stand still or jump

- Let’s be more careful (and ambitious):
- Logical clocks should always move forward
  - Sometimes faster, sometimes slower is OK.
  - But there should be a minimum and a maximum speed.
  - As close to correct time as possible!

Formal Model

- Hardware clock $H_v(t) = \int_{[0,t]} h_v(\tau)\, d\tau$ with clock rate $h_v(t) \in [1-\epsilon, 1+\epsilon]$

Clock drift ϵ is typically small, e.g. $\epsilon \approx 10^{-4}$ for a cheap quartz oscillator

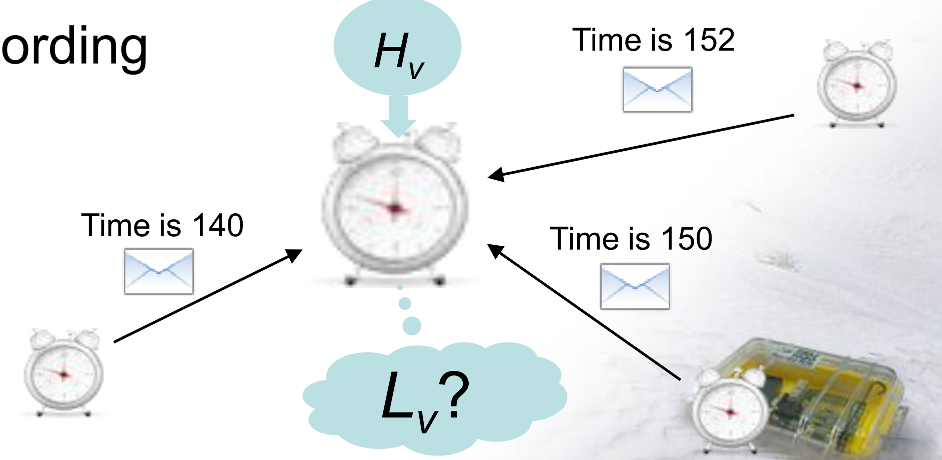
- Logical clock $L_v(\cdot)$ which increases at rate at least 1 and at most β

Logical clocks with rate less than 1 behave differently ("synchronizer")

- Message delays $\in [0,1]$

Neglect fixed share of delay, normalize jitter

- Employ a synchronization algorithm to update the logical clock according to hardware clock and messages from neighbors

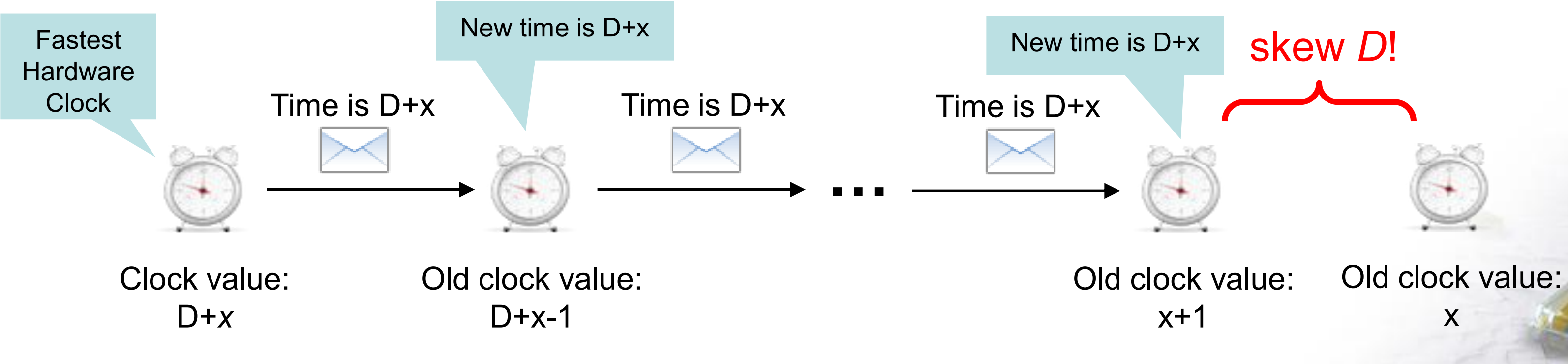


Synchronization Algorithms: An Example (“$A^{\max}$”)

- Question: How to update the logical clock based on the messages from the neighbors?

Allow $\beta = \infty$

- Idea: Minimizing the skew to the fastest neighbor
  - Set the clock to the maximum clock value received from any neighbor (if larger than local clock value)
  - forward new values immediately
- Optimum global skew of about *D*
- Poor local property
  - First all messages take 1 time unit…
  - …then we have a fast message!



Fastest Hardware Clock

Time is D+x

New time is D+x

Time is D+x

…

Time is D+x

New time is D+x

skew *D*!

Clock value: D+*x*

Old clock value: D+x-1

Old clock value: x+1

Old clock value: x

Synchronization Algorithms: A

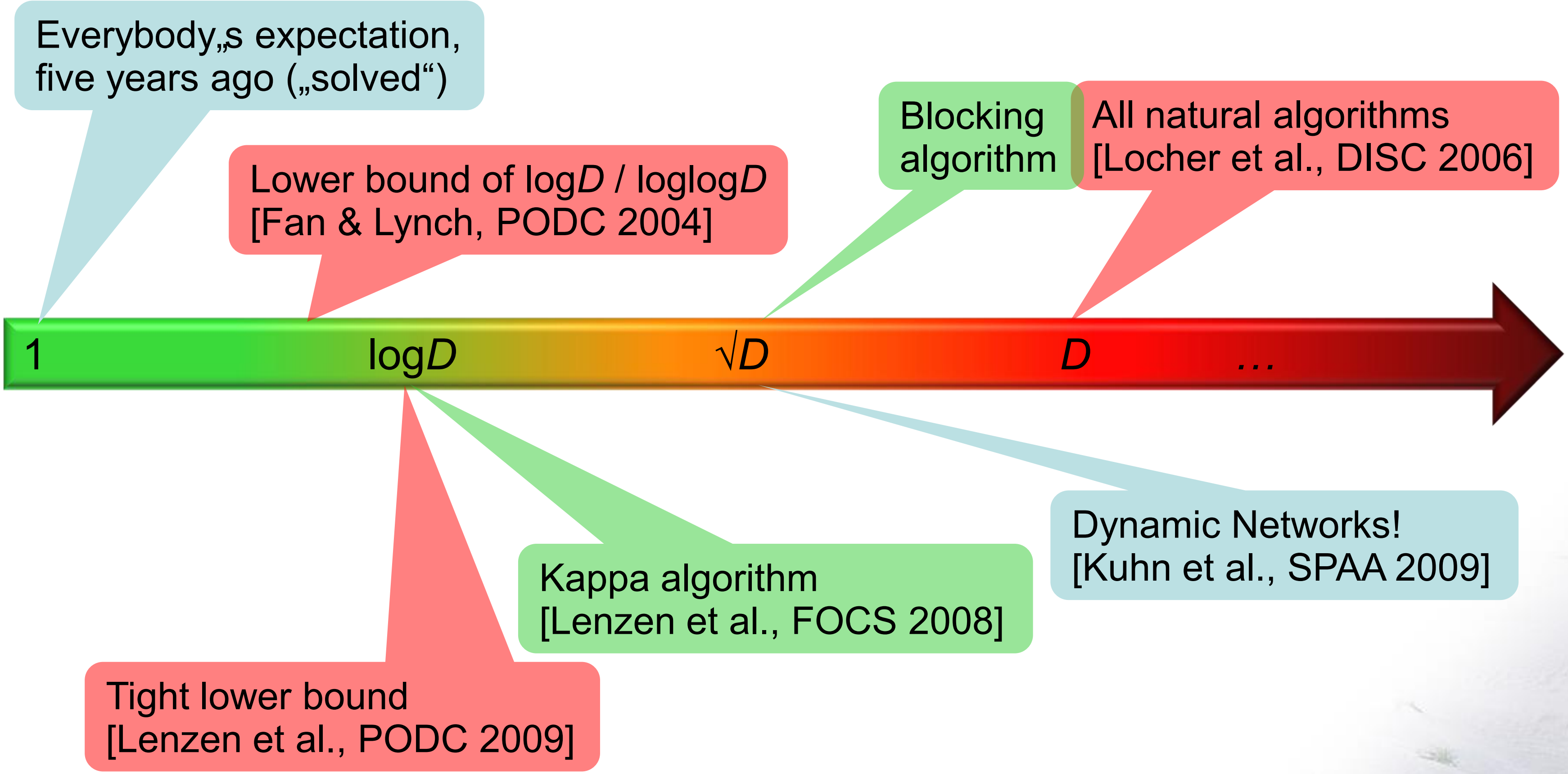
- The problem of A^{max} is that the cl
  maximum value
- Idea: Allow a constant slack γ be
  value and the own clock value
- The algorithm $A^{max'}$ sets the local

$$L_i(t) := \max(\, L_i(t), \max_{j \in N_i} L_j(t) - \gamma\,)$$

→ Worst-case clock skew between two neighboring nodes is still Θ(D) independent of the choice of γ!

- How can we do better?
  - Adjust logical clock speeds to catch up with fastest node (i.e. no jump)?
  - Idea: Take the clock of all neighbors into account by choosing the average value?

Local Skew: Overview of Results



Everybody„s expectation, five years ago („solved“)

Lower bound of log$D$ / loglog$D$
[Fan & Lynch, PODC 2004]

Blocking algorithm

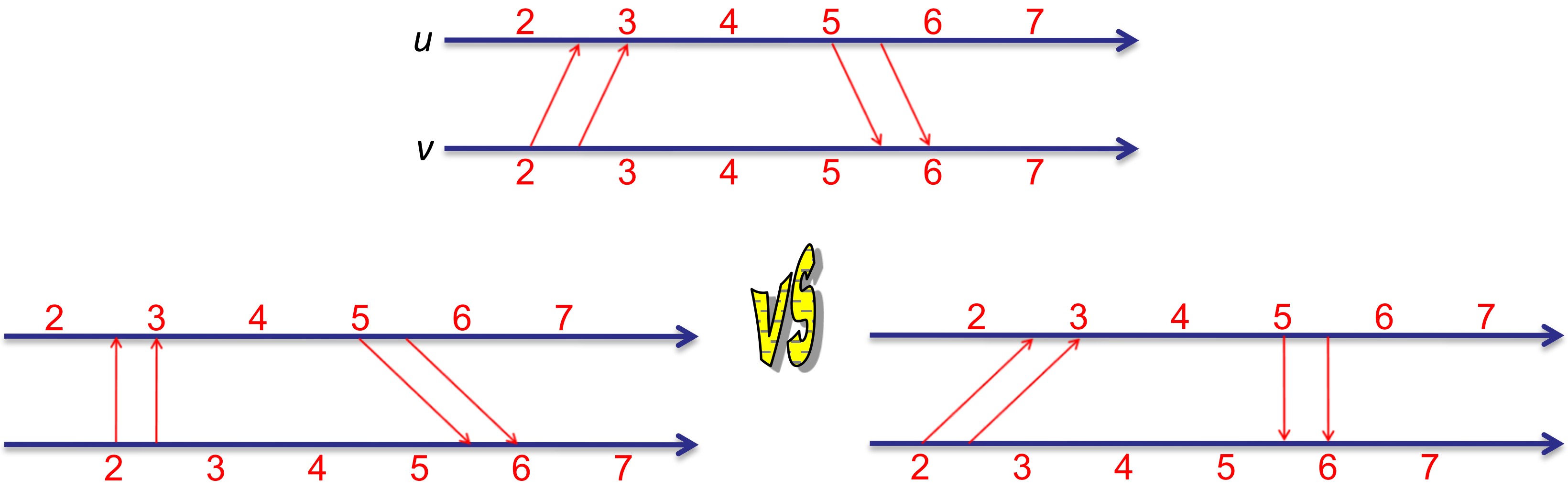
All natural algorithms
[Locher et al., DISC 2006]

Tight lower bound
[Lenzen et al., PODC 2009]

Kappa algorithm
[Lenzen et al., FOCS 2008]

Dynamic Networks!
[Kuhn et al., SPAA 2009]

Enforcing Clock Skew

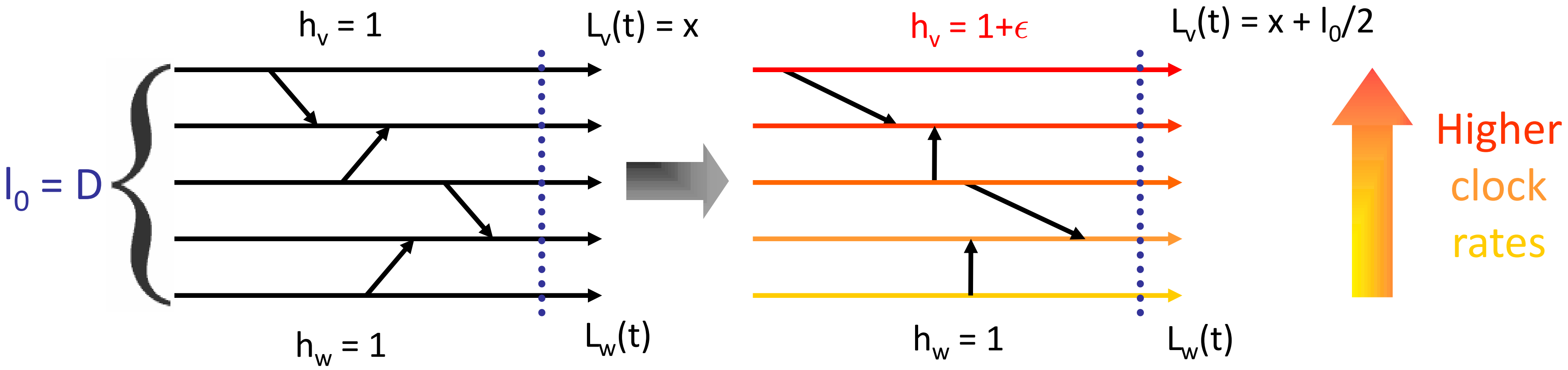


- Messages between two neighboring nodes may be fast in one direction and slow in the other, or vice versa.

- A constant skew between neighbors may be „hidden“.

- In a path, the global skew may be in the order of *D*/2.

Local Skew: Lower Bound (Single-Slide Proof!)



- Add $l_0/2$ skew in $l_0/(2\epsilon)$ time, messing with clock rates and messages
- Afterwards: Continue execution for $l_0/(4(\beta-1))$ time (all $h_x = 1$)
  - → Skew reduces by at most $l_0/4$ → at least $l_0/4$ skew remains
  - → Consider a subpath of length $l_1 = l_0 \cdot \epsilon/(2(\beta-1))$ with at least $l_1/4$ skew
  - → Add $l_1/2$ skew in $l_1/(2\epsilon) = l_0/(4(\beta-1))$ time → at least $3/4 \cdot l_1$ skew in subpath
- Repeat this trick (+½,-¼,+½,-¼,…) $\log_{2(\beta-1)/\epsilon} D$ times

Theorem: $\Omega(\log_{(\beta-1)/\epsilon} D)$ skew between neighbors

Local Skew: Upper Bound

- Surprisingly, up to small constants, the $\Omega(\log_{(\beta-1)/\epsilon} D)$ lower bound can be matched with clock rates $\in [1,\beta]$
- We get the following picture [Lenzen et al., PODC 2009]:

max rate β	$1+\epsilon$	$1+\Theta(\epsilon)$	$1+\sqrt{\epsilon}$	2	large
local skew	∞	$\Theta(\log D)$	$\Theta(\log_{1/\epsilon} D)$	$\Theta(\log_{1/\epsilon} D)$	$\Theta(\log_{1/\epsilon} D)$

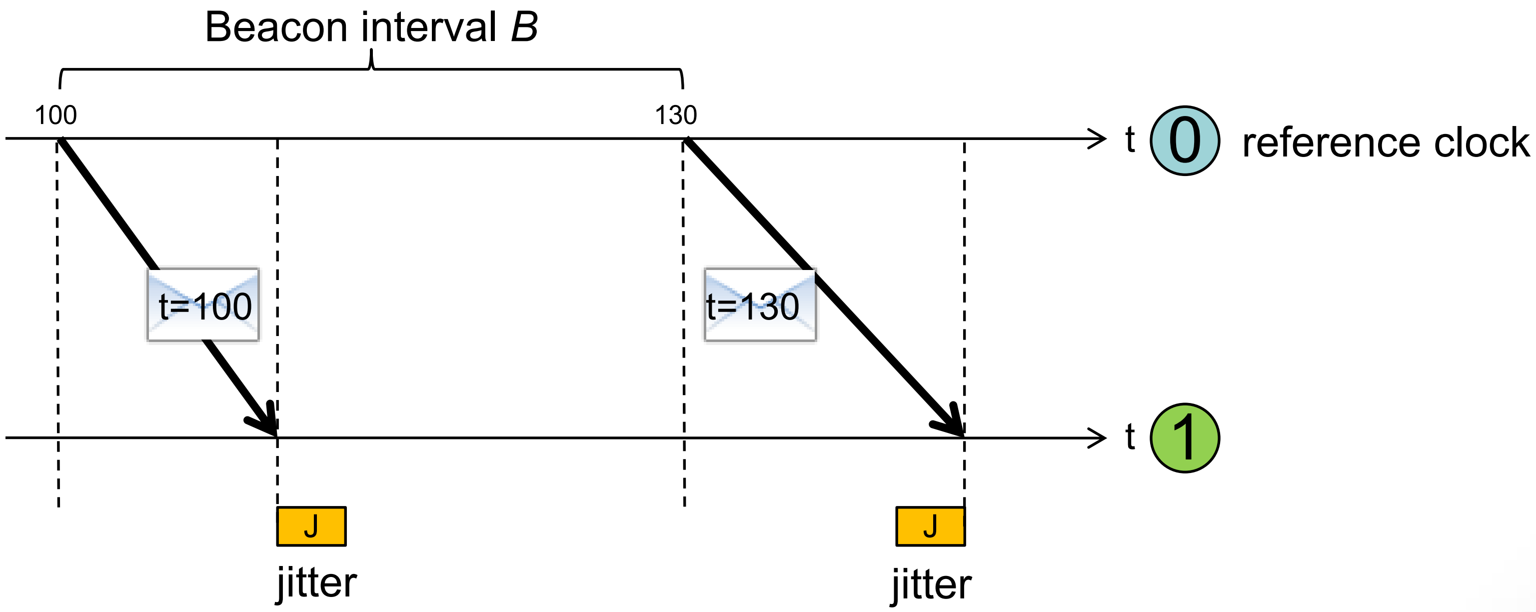
We can have both smooth and accurate clocks!

... because too large clock rates will amplify the clock drift ϵ.

- In practice, we usually have $1/\epsilon \approx 10^4 > D$. In other words, our initial intuition of a constant local skew was not entirely wrong! ☺

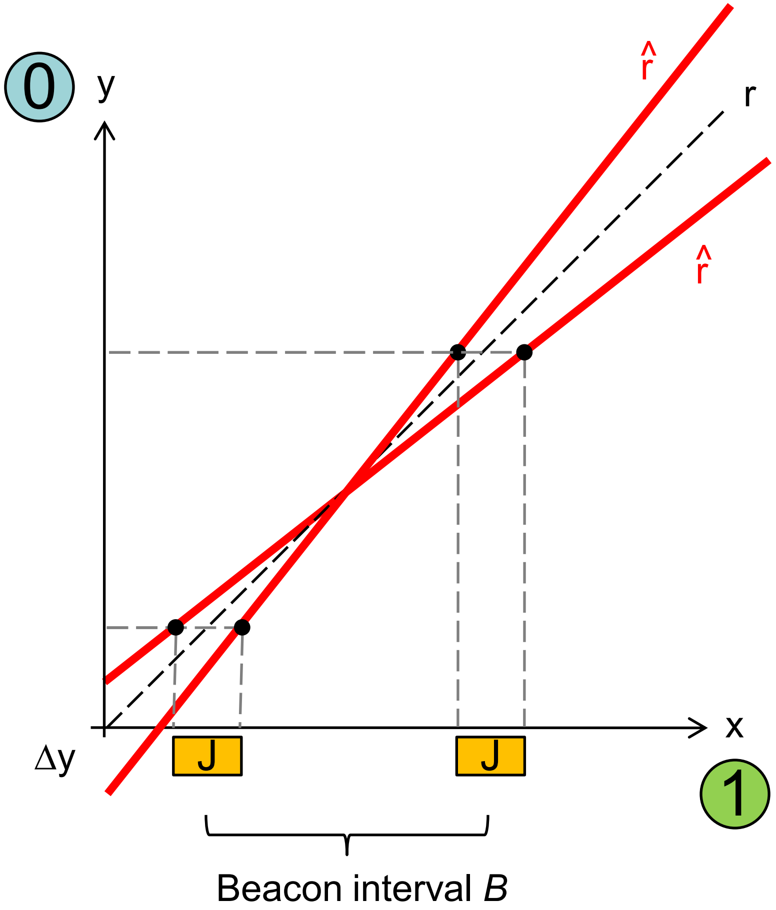
Synchronizing Nodes

- Sending periodic beacon messages to synchronize nodes



How accurately can we synchronize two Nodes?

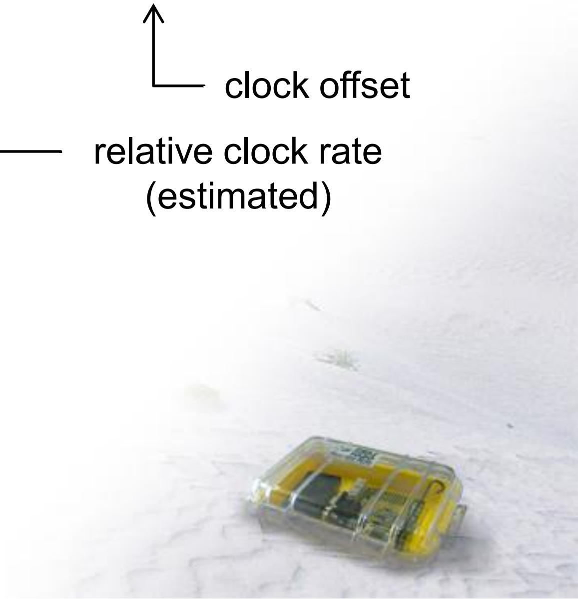
- Message delay jitter affects clock synchronization quality



$$y(x) = \hat{r} \cdot x + \Delta y$$

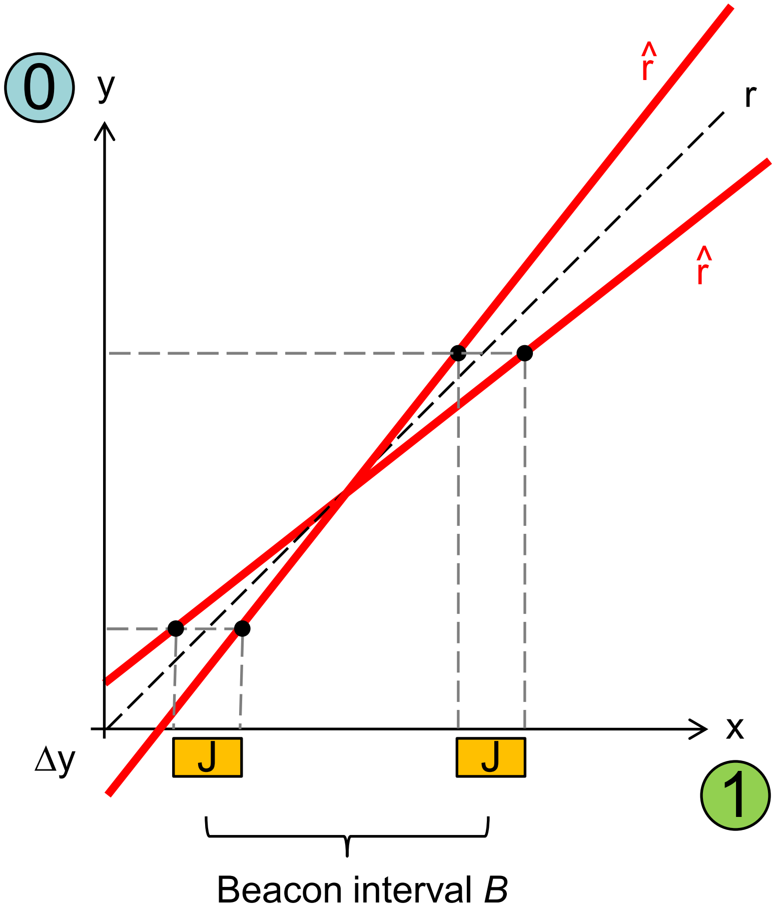
clock offset

relative clock rate
(estimated)



Clock Skew between two Nodes

- Lower Bound on the clock skew between two neighbors



Error in the rate estimation:
– Jitter in the message delay
– Beacon interval
– Number of beacons k

$$|\hat{r} - r| \sim \frac{J}{Bk\sqrt{k}}$$

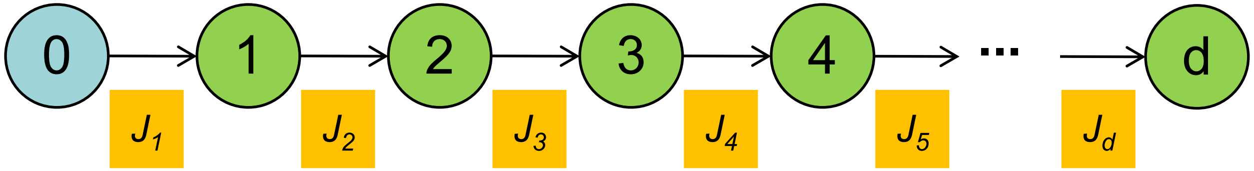
Synchronization error:

$$|\hat{y} - y| \sim \frac{J}{\sqrt{k}}$$

Multi-hop Clock Synchronization

- Nodes forward their current estimate of the reference clock

  Each synchronization beacon is affected by a **random jitter *J***



- Sum of the jitter grows with the square-root of the distance

  $\mathrm{stddev}(J_1 + J_2 + J_3 + J_4 + J_5 + \ldots J_d) = \sqrt{d} \times \mathrm{stddev}(J)$

Single-hop:

$$|\hat{y} - y| \sim \frac{J}{\sqrt{k}}$$

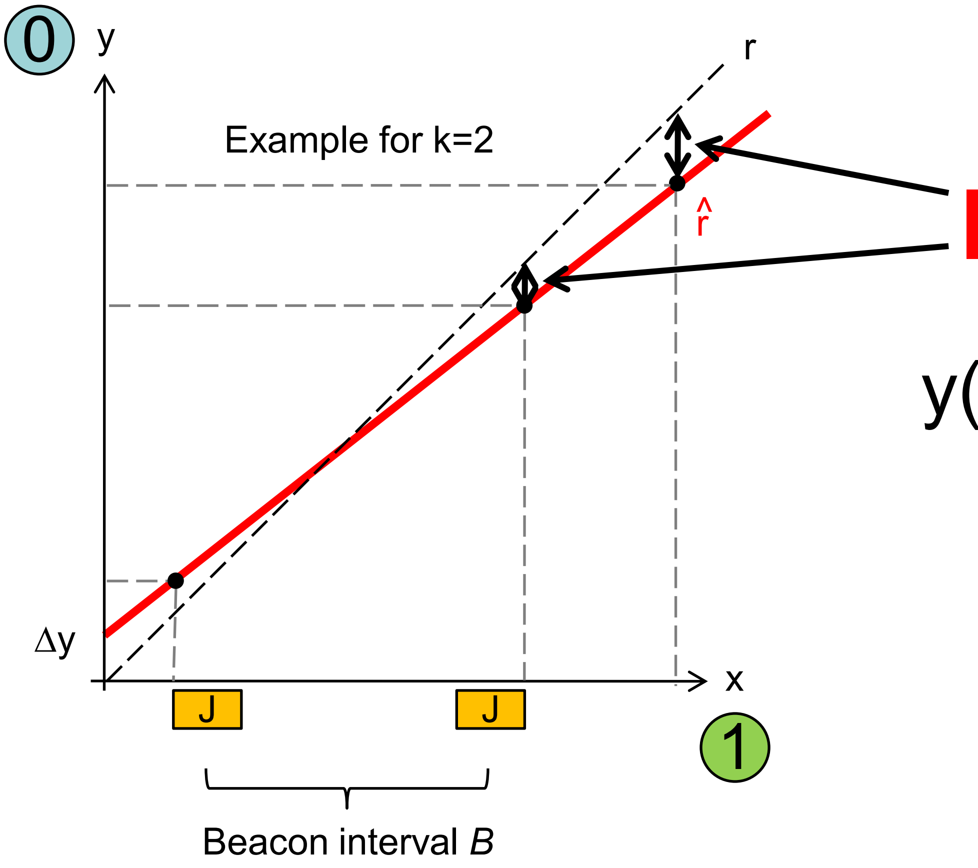
$\Rightarrow$

Multi-hop:

$$|\hat{y} - y| \sim \frac{J\sqrt{d}}{\sqrt{k}}$$

Linear Regression (e.g. FTSP)

- FTSP uses linear regression to compensate for clock drift
  Jitter is amplified before it is sent to the next hop



synchronization error

$$y(x) = \hat{r} \cdot x + \Delta y$$

clock offset

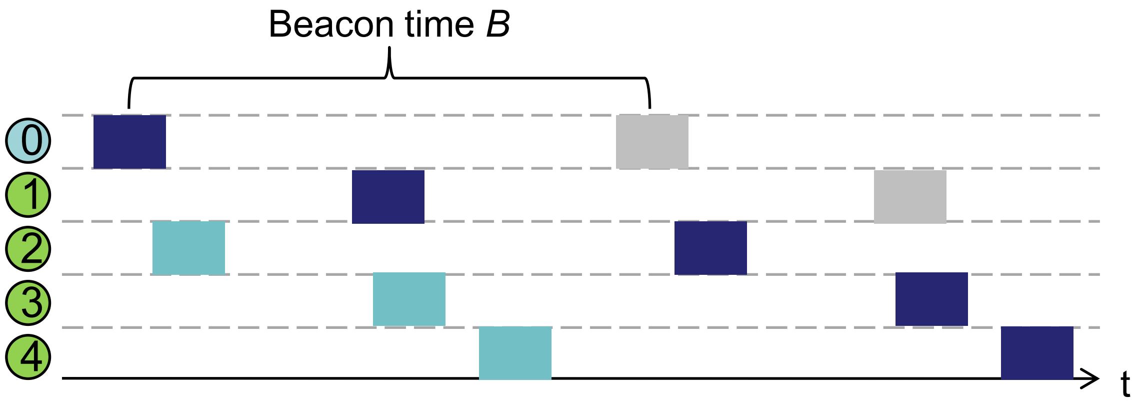
relative clock rate (estimated)

The PulseSync Protocol

- Send fast synchronization pulses through the network
  - Speed-up the initialization phase
  - Faster adaptation to changes in temperature or network topology

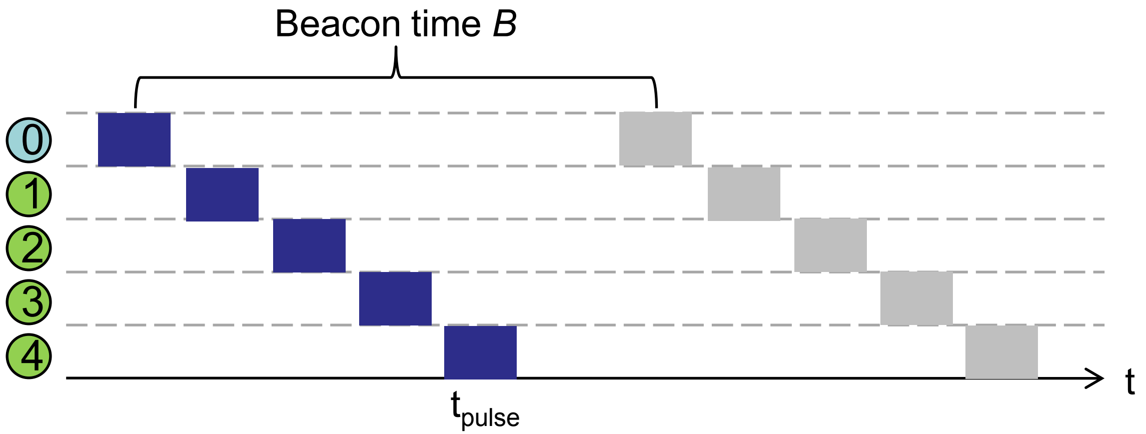
FTSP

Expected time
= $D \cdot B/2$



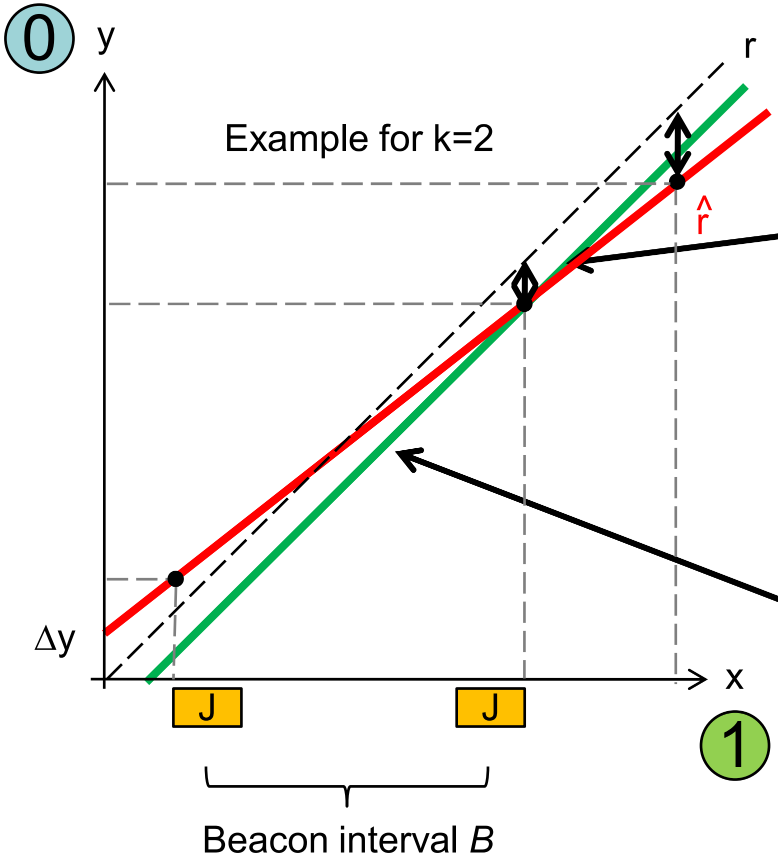
PulseSync

Expected time
= $D \cdot t_{pulse}$



The PulseSync Protocol (2)

- Remove self-amplification of synchronization error
  - Fast flooding cannot completely eliminate amplification



synchronization error

$$y(x) = \hat{r} \cdot x + \Delta y$$

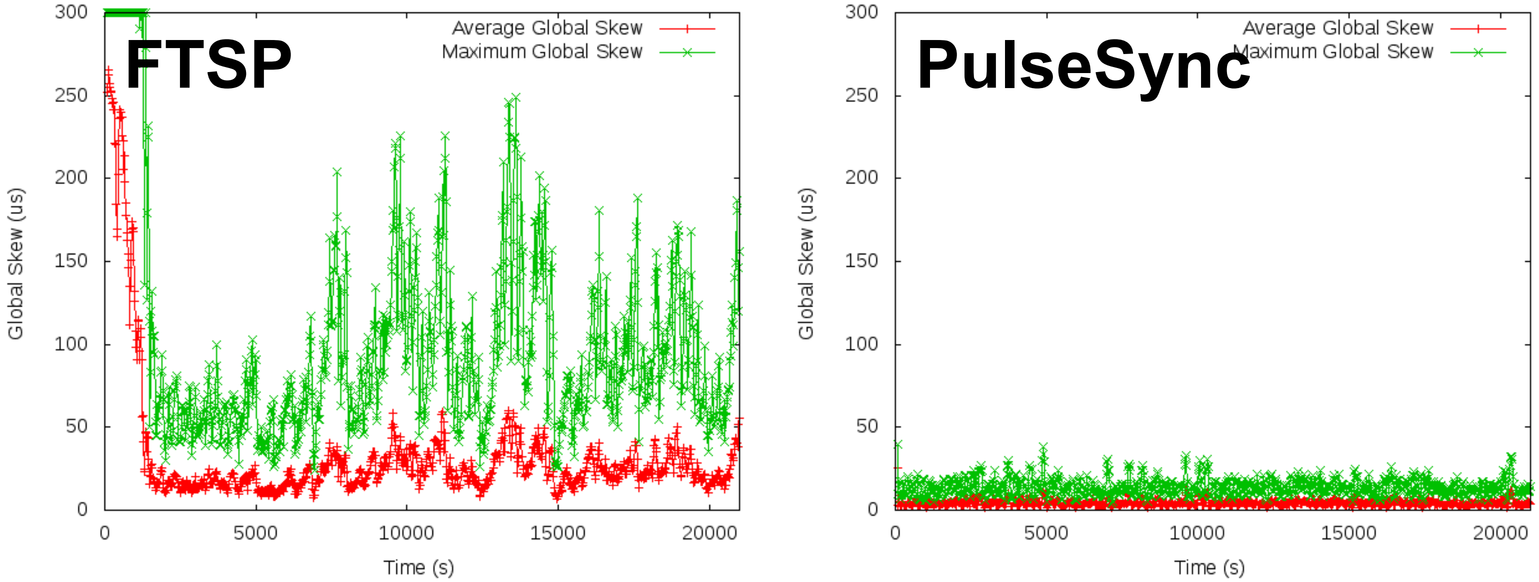
clock offset

relative clock rate
(estimated)

The green line is calculated using
k measurement points that are
statistically *independent* of the red line.

FTSP vs. PulseSync

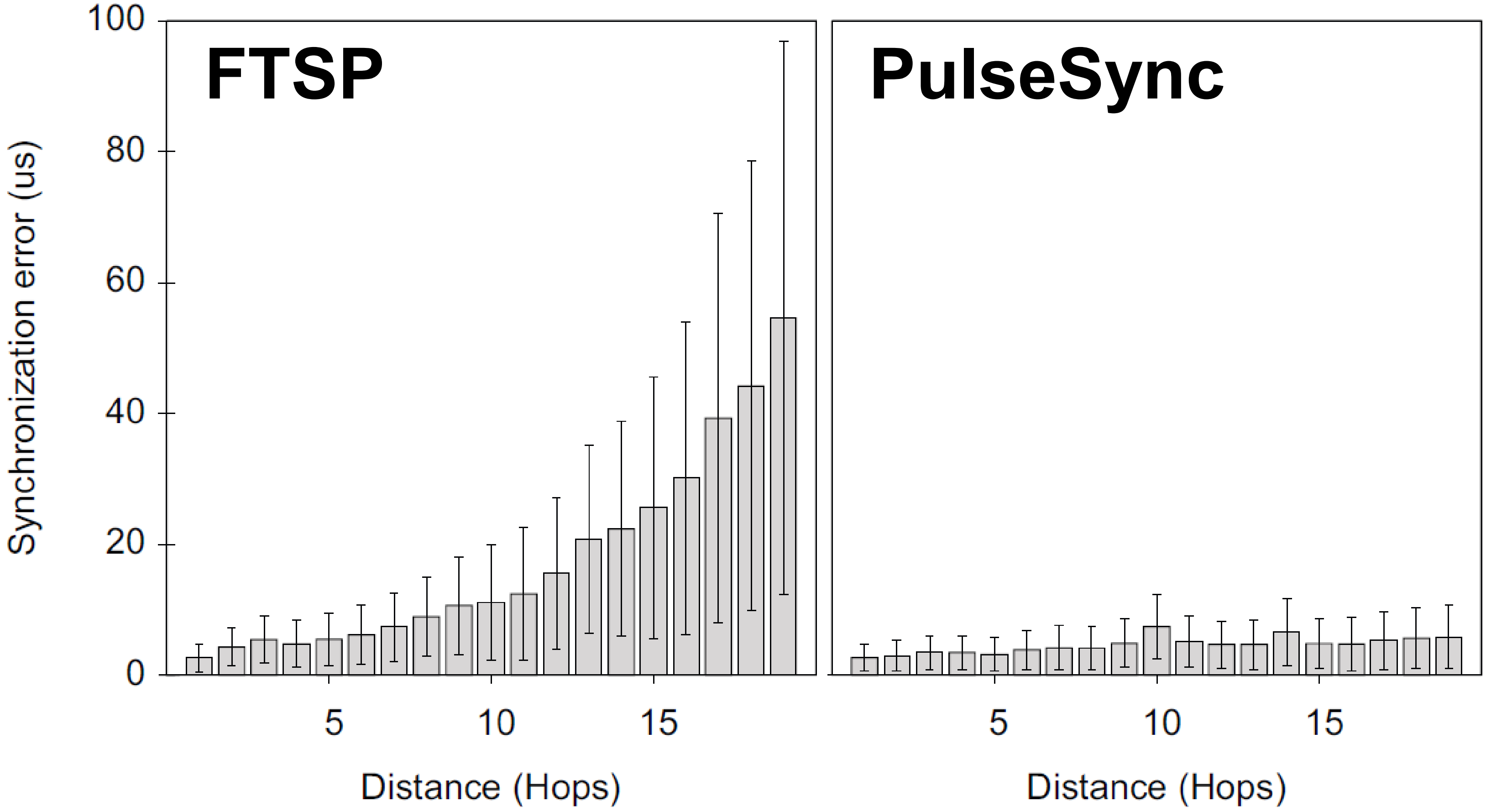
- Global Clock Skew
  - Maximum synchronization error between any two nodes



Synchronization Error	FTSP	PulseSync
Average (t>2000s)	23.96 μs	4.44 μs
Maximum (t>2000s)	249 μs	38 μs

FTSP vs. PulseSync

- Sychnronization Error vs. distance from root node



Open Problem

- As listed on slide 9/6, clock synchronization has lots of parameters. Some of them (like local/gradient) clock synchronization have only started to be understood.

- Local clock synchronization in combination with other parameters are not understood well, e.g.
  - accuracy vs. convergence
  - fault-tolerance in case some clocks are misbehaving [Byzantine]
  - clock synchronization in dynamic networks