



# Principles of Distributed Computing

## Exercise 1

### 1 Vertex Coloring

- In the lecture, a simple distributed algorithm which colors an arbitrary graph with  $\Delta + 1$  colors in  $n$  synchronous rounds was presented ( $\Delta$  denotes the greatest degree,  $n$  the number of nodes of the graph). To run in  $n$  rounds, the nodes of the graph had to be numbered from 1 to  $n$ . Devise a synchronous distributed algorithm for the case the IDs are unique but unbounded numbers (i.e. the nodes have arbitrary IDs instead of being numbered from 1 to  $n$ ). Your algorithm should also use at most  $\Delta + 1$  colors and terminate in a linear number of synchronous rounds.
- What is the total number of messages your algorithm sends?
- Does your algorithm also work in an asynchronous environment? If yes, formulate the asynchronous equivalent to your algorithm, if no, describe why.

### 2 Counting the Nodes of a Tree

In this exercise, we assume that the communication graph  $T$  is a tree. We consider different aspects of the problem of counting the number of nodes of  $T$ .

- Suppose that a node  $v \in T$  wants to know the total number of nodes. Develop a distributed algorithm  $\mathcal{A}$  for this task.  $\mathcal{A}$  can be started by every node  $v$  of  $T$ , it should determine the number of nodes of  $T$  and report it to  $v$ . How long does your algorithm need until  $v$  knows the result?
- Suppose now that all nodes would like to know the number of nodes in  $T$ . Devise an algorithm with which all nodes of a tree  $T$  concurrently calculate the number of nodes.
- For the last question of this exercise, we assume that the tree  $T$  has an odd number of nodes. In such a tree, there is a unique node  $v$  which allows to divide  $T$  into two parts whose sizes are as equal as possible. Can you use your results of Question 2b) to develop a distributed algorithm to find  $v$ ? What can you say about the sizes of both parts of the achieved partition of  $T$ ?