



# Principles of Distributed Computing

## Exercise 9

### 1 Pancake Networks

In the lecture, you have encountered several different graphs as underlying network structures for peer-to-peer (P2P) networks. Here, we will look at another prominent example, the *Pancake graph*<sup>1</sup>  $P_n$ .

The pancake graph  $P_n$  is defined as follows: The vertex set is

$$V(P_n) = \{(v_1, v_2, \dots, v_n) \mid v_i \in \{1, \dots, n\} \text{ and } v_i \neq v_j \forall i \neq j\}.$$

In other words,  $V(P_n) = S_n$ , the group of all permutations on  $n$  elements. There exists an edge of dimension  $i$  for  $2 \leq i \leq n$  when

$$e_i = \{(v_1, \dots, v_i \dots, v_n), (w_1, \dots, w_i, \dots, w_n)\} \in E(P_n) \iff w_j = v_{i-j+1} \text{ for } 1 \leq j \leq i \text{ and } w_j = v_j \text{ for } i < j \leq n,$$

or, we can say that an edge  $e_i$  represents a *prefix reversal*

$$(v_1, \dots, v_i, v_{i+1}, \dots, v_n) \longleftrightarrow (v_i, \dots, v_1, v_{i+1}, \dots, v_n). \quad (1)$$

For the following questions, where appropriate, give your answers in terms of  $N := |V(P_n)|$  (approximately), the number of vertices, as well as  $n$ .

- Draw (nicely!)  $P_n$  for  $n = 2, 3, 4$ . Try to describe a pattern for drawing  $P_n$  for any  $n$ .
- What is the degree of each vertex in  $P_n$ ?
- Can you give bounds on the diameter  $D(P_n)$  of the pancake network?
- Show that  $P_n$  is Hamiltonian for  $n \geq 3$ .
- How can the pancake graph be used to implement a *distributed hash table* (DHT)? In other words, where are files, indexed by bitstrings of a certain length  $b$ , stored in the pancake graph, and how can these files be looked up (given the corresponding bitstring)?<sup>2</sup>

The pancake graph has been proposed for P2P networks partly because of the properties analyzed in this exercise.

<sup>1</sup>A well-known paper about pancake graphs was originally written in 1976 by a 21 year old college dropout (later to become a famous entrepreneur) and a Ph.D. student (later to become a famous scholar):  
W. Gates and C. Papadimitriou. Bounds for Sorting by Prefix Reversal. *Discrete Math.*, 27:4757, 1979.

<sup>2</sup>You can ignore churn in this exercise.